# **Finding Global Minima via Kernel Approximations**

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Joint work with Alessandro Rudi and Francis Bach LAAS - February 16, 2020

## Outline

- Global optimization using function values
- Representing non-negative functions
- Global optimization through sampling
- Conclusions and extensions

## **Global optimization using function values**

- Global optimization
- Optimal algorithms
- Optimal rates
- Convex reformulation

## **Global optimization**

• Zero-th order minimization

 $\min_{x \in \Omega} f(x)$ 

- $\Omega \subset \mathbb{R}^d$  simple compact subset (e.g.,  $[-1,1]^d$ )
- -f with some bounded derivatives
- access to function values

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- -f with some bounded derivatives
- access to function values
- No convexity assumption
- Many applications
  - e.g., hyperparameter optimization in machine learning

## **Optimal algorithms**

- Goal: Find  $\hat{x} \in \Omega$  such that  $f(\hat{x}) \min_{x \in \Omega} f(x) \leqslant \varepsilon$ 
  - Lowest number of function calls
  - Worst-case guarantees over all functions f in some convex set  ${\mathcal F}$

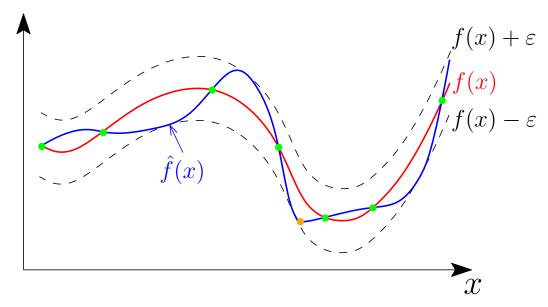
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- Equivalence to uniform function approximation (Novak, 2006)
  - Simplest algorithm: approximate f by  $\hat{f}$  and minimize  $\hat{f}$



### **Optimal rates**

- Optimal worst-case performance over  $\mathcal{F}$  (Novak, 2006)
  - -n = number of function evaluations
  - $\mathcal{F} = \text{Lipschitz-continuous functions:} n \propto \varepsilon^{-d}$
  - $\mathcal{F} = m$  bounded derivatives:  $n \propto \varepsilon^{-d/m}$
- Smoothness to circumvent the curse of dimensionality
  - NB: constants may depend (exponentially) in d

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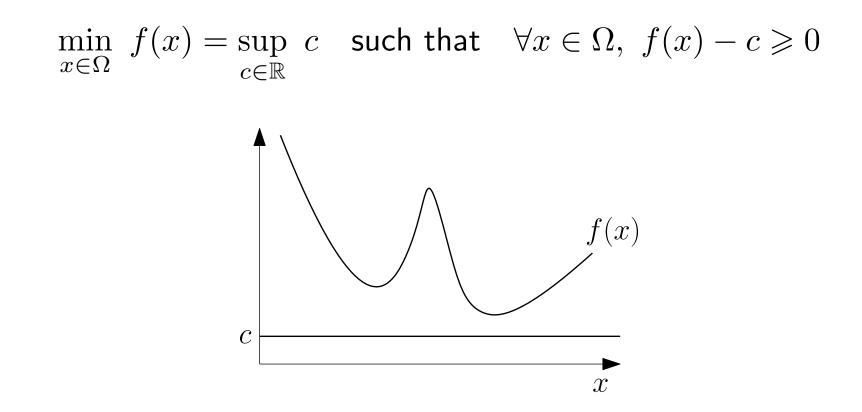
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  - "Approximate then optimize"
- Algorithms with polynomial-time complexity in n?
  - "Approximate and optimize"

#### Reformulations

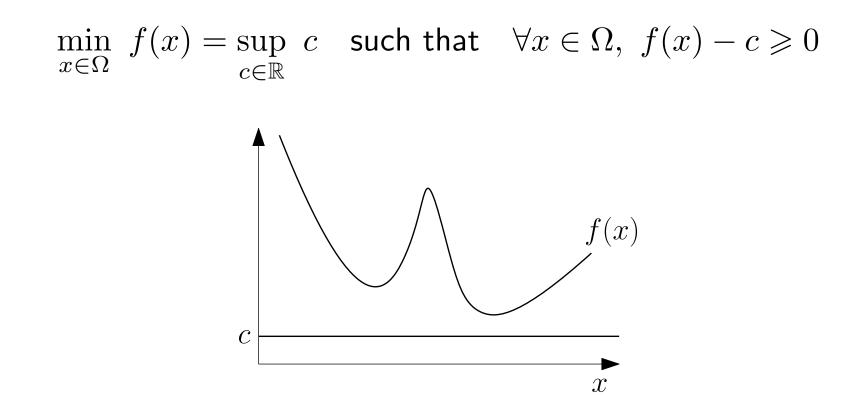
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#### Reformulations

#### • Equivalent convex problem



- All optimization problems are convex!
- Need to represent non-negative functions (such as f(x) c)

## **Representing non-negative functions**

- Generic representation in a PSD cone
- The case of polynomials
- Infinite-dimensional representations using RKHS

### **Representing non-negative functions**

- Assumption: g(x) can be represented as  $g(x) = \langle \phi(x), G\phi(x) \rangle$  with G symmetric operator
- **Example**: set of polynomials of degree 2rwith  $\phi(x)$  composed of monomials of degree r, of dimension  $\binom{d+r}{r}$

$$\phi(x) = (1, x_1, \dots, x_d, x_1 x_2, \dots, x_1^r, \dots, x_d^r)$$

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- Non-negativity through "sums-of-squares"
  - If  $G \succcurlyeq 0$ , then  $\forall x \in \Omega, \ g(x) = \langle \phi(x), G \phi(x) \rangle \ge 0$
  - Sum of squares : eigen-decomposition  $G = \sum_{i \in I} \lambda_i \ h_i \otimes h_i$ ,

$$g(x) = \sum_{i \in I} \lambda_i \langle \phi(x), (h_i \otimes h_i) \phi(x) \rangle = \sum_{i \in I} \lambda_i \langle \phi(x), h_i \rangle^2$$

### **Global optimization with sums of square polynomials**

- Replace  $f c \ge 0$  by  $f c = \langle \phi(x), A \phi(x) \rangle$  with  $A \succcurlyeq 0$
- **Sum-of-squares optimization** (Lasserre, 2001)

 $\sup_{c \in \mathbb{R}, A \succeq 0} c \quad \text{such that} \quad \forall x \in \mathbb{R}^d, \ f - c = \langle \phi(x), A \phi(x) \rangle \quad (\mathsf{P})$ 

- Equivalent to original problem if  $f f_*$  is SoS
- Polynomial constraints can be added (and help !)
- If not sum of squares : **hierarchies**  $(P_i)$  of degree  $i \to \infty$

#### • Guarantees

- the hierarchy converges to the optimum value
- Lower bound when solving  $P_i$ :  $f_i \leq f_*$ .
- Solve a finite dimensional SDP

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- If not sum of squares : **hierarchies**  $(P_i)$  of degree  $i \to \infty$

#### • Drawbacks

- Not all non-negative polynomials are SoS if  $d\geq 1$
- No guarantee a priori on the degree i needed for a given precision
- Only for polynomials

#### **Representing more general functions with RKHS**

- Reproducing Kernel Hilbert Space (RKHS) :
  - Hilbert space of functions  $g\in\mathcal{H},\ g:\mathbb{R}^d\rightarrow\mathbb{R}$
  - Reproducing property :  $g(x) = \langle g, \phi(x) \rangle_{\mathcal{H}}$
  - Kernel :  $k(x, x') = \langle \phi(x), \phi(x') \rangle_{\mathcal{H}}$  (computable)

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- Example : Sobolev spaces (Berlinet and Thomas-Agnan, 2011)
  - Sobolev spaces  $H^s(\Omega)$  with  $\Omega \subset \mathbb{R}^d$ , s > d/2

$$\langle f,g\rangle_{H^s(\Omega)} = \sum_{|\alpha| \leq s} \int_{\Omega} \partial^{\alpha} f \ \partial^{\alpha} g$$

$$-s = d/2 + 1/2 : k(x, y) = \exp(-\|x - y\|)$$

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#### $\bullet$ Everything can be expressed using only the kernel function k

- k is known for all Sobolev spaces with s > d/2.
- $-s = d/2 + 1/2 : k(x, y) = \exp(-\|x y\|)$

#### **Non-negative functions with RKHS**

- **RKHS** *H* :
  - Feature map  $\phi(x) \in \mathcal{H}$   $(g(x) = \langle g, \phi(x) \rangle_{\mathcal{H}})$
  - Kernel :  $k(x, x') = \langle \phi(x), \phi(x') \rangle_{\mathcal{H}}$  (computable)
- RKHS Sum of Squares : (Marteau-Ferey et al., 2020)
  - $-g_A(x) = \langle \phi(x), A\phi(x) \rangle, \quad A \succeq 0$
  - Eigendecomposition  $A = \sum_{i \in I} \lambda_i h_i \otimes h_i$ SoS :  $g_A(x) = \sum_{i \in I} \lambda_i h_i(x)^2, h_i \in \mathcal{H}$

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- $\bullet$  Everything can be expressed using the kernel k
  - Sum of squares of sobolev functions.

## **Global optimization through sampling**

- Global optimization : formulation and theoretical result in the RKHS
- Controlled approximation through sampling
- Algorithms and illustrations

## Going infinite-dimensional (Rudi, Marteau-Ferey, and Bach, 2020)

 $\sup_{c\in\mathbb{R},\ A\succcurlyeq 0} c \quad \text{such that} \quad \forall x\in\Omega,\ f(x)-c=\langle \phi(x),A\phi(x)\rangle$ 

- Sobolev space  $\mathcal{H}$ :
  - -s > d/2 squared-integrable derivative
  - RKHS with feature map  $\phi(x)$  and kernel k(x, x')
  - RKHS SoS :  $g_A(x) = \langle \phi(x), A\phi(x) \rangle_{\mathcal{H}}, \ A \succeq 0$

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- Theorem:  $\exists A_* \geq 0$  such that  $\forall x \in \Omega$ ,  $f(x) = f_* + \langle \phi(x), A_*\phi(x) \rangle$ 
  - If f has isolated strict-second order minima in  $\check{\Omega},$  and f is (s+3)-times differentiable

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  - If f has isolated strict-second order minima in  $\check{\Omega},$  and f is (s+3)-times differentiable
  - $\Rightarrow$  Equivalent to original problem, but infinite-dimensional

• Subsample n points  $x_1, \ldots, x_n \in \Omega$  and solve

 $\sup_{c \in \mathbb{R}, A \succeq 0} c - \lambda \operatorname{tr}(A) \text{ such that } \forall i \in \{1, \dots, n\}, f(x_i) = c + \langle \phi(x_i), A \phi(x_i) \rangle$ 

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- Finite-dimensional formulation through representer theorem (kernels !)
  - Marteau-Ferey, Bach, and Rudi (2020)
  - SDP of dimension n :

 $\sup_{c \in \mathbb{R}, \ B \succeq 0, B \in \mathbb{R}^{n \times n}} c - \lambda \operatorname{tr}(B) \quad \text{st } \forall i \in \{1, \dots, n\}, \ f(x_i) = c + \Phi_i^\top B \Phi_i$ 

• Solvable in polynomial time with precision  $\epsilon$  in  $O(n^{3.5} \log \frac{1}{\epsilon})$ 

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- **Theorem** (Rudi, Marteau-Ferey, and Bach, 2020) Let  $\hat{c}, \hat{A}$  be the result of the algorithm. Up to logarithmic terms :
  - if  $n = \Omega(\varepsilon^{-d/(m-d/2-3)})$  and the samples  $(x_1, ..., x_n)$  are taken randomly from  $\Omega$ , and if  $\lambda = \varepsilon$ , then it holds with probability at least  $1 \delta$ :

$$|\hat{c} - f_*| \le \varepsilon \log \frac{1}{\delta}$$

- in that case, the complexity is  $\Omega(\varepsilon^{-3d/(m-d/2-3)}\log \frac{1}{\epsilon})$ 

• **Extension**: possible to find the minimizer  $\hat{x}$ .

• Subsample n points  $x_1, \ldots, x_n \in \Omega$  and solve

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- Approximation guarantees (Rudi, Marteau-Ferey, and Bach, 2020)
  - With random samples,  $n \approx \varepsilon^{-d/(m-d/2-3)}$ 
    - (up to logarithmic terms)
  - To be compared to optimal rate  $n\approx \varepsilon^{-d/(m-d/2)}$
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- Subsampling inequalities as  $f(x_i) \ge c$  directly?
  - cannot improve on  $n \approx \varepsilon^{-d}$

## **Final algorithm**

- Input:  $f: \mathbb{R}^d \to \mathbb{R}$ ,  $\Omega \subset \mathbb{R}^d, n \ge 0, \lambda > 0, s > d/2$
- 1. Sampling:  $\{x_1, \ldots, x_n\}$  sampled i.i.d. uniformly on  $\Omega$

#### 2. Feature computation

- Set  $f_j = f(x_j)$ ,  $\forall j \in \{1, \ldots, n\}$ 

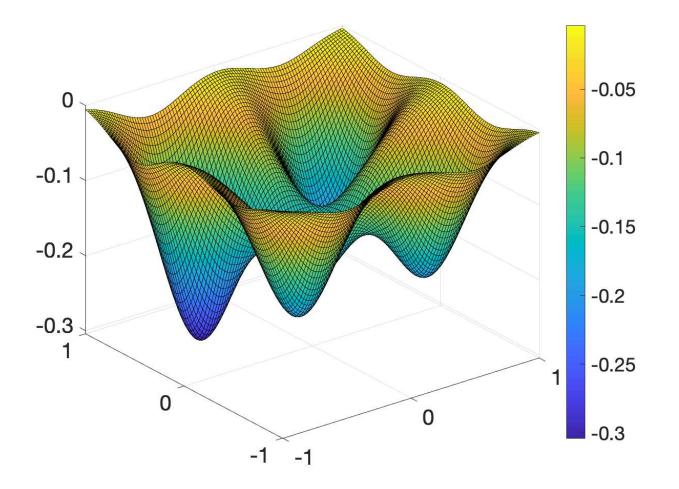
- Compute  $K_{ij} = k(x_i, x_j)$  for k Sobolev kernel of smoothness s
- Set  $\Phi_j \in \mathbb{R}^n$  computed using a cholesky decomposition of K $\forall j \in \{1, \dots, n\}.$
- 3. Solve  $\max_{c \in \mathbb{R}, B \succeq 0} c \lambda \operatorname{tr}(B)$  s.t.  $\forall j \in \{1, \dots, n\}, f_j c = \Phi_j^\top B \Phi_j$
- **Output:** c proxy for  $f_*$
- One can extend the algorithm in order to compute a proxy of the minimizer

## **Opposite properties from Polynomial SoS !**

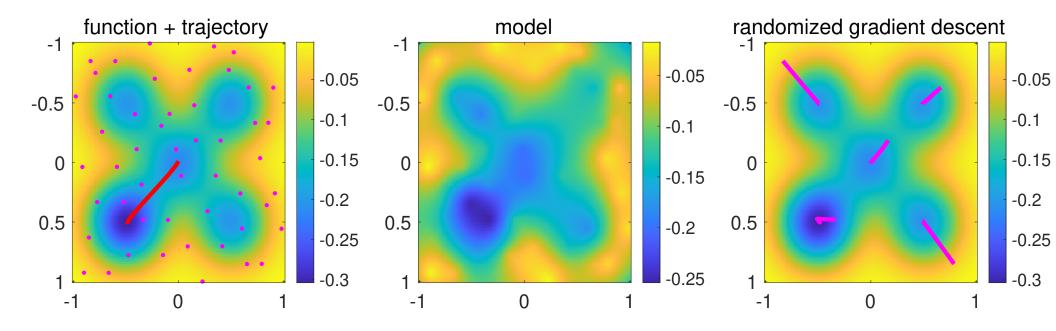
- "Always" possible to write a non-negative function as a RKHS SoS
- Bounds on the number of samples needed for a given precision
- Finite dimensional SDP with bounded complexity  $O(n^{3.5} \log \frac{1}{\epsilon})$
- Breaks the curse of dimensionality (needs  $\epsilon^{-d/m}$  samples) for smooth enough functions
- For the moment, no certificate bound on the result of the algorithm

#### Illustration

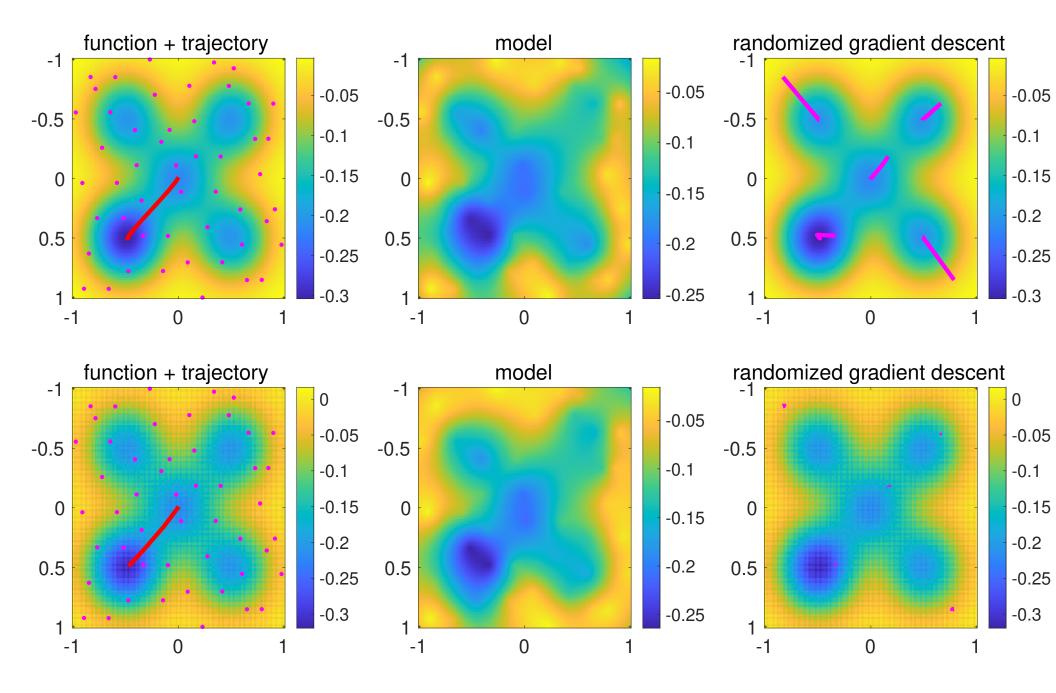
• Minimization of two-dimensional function



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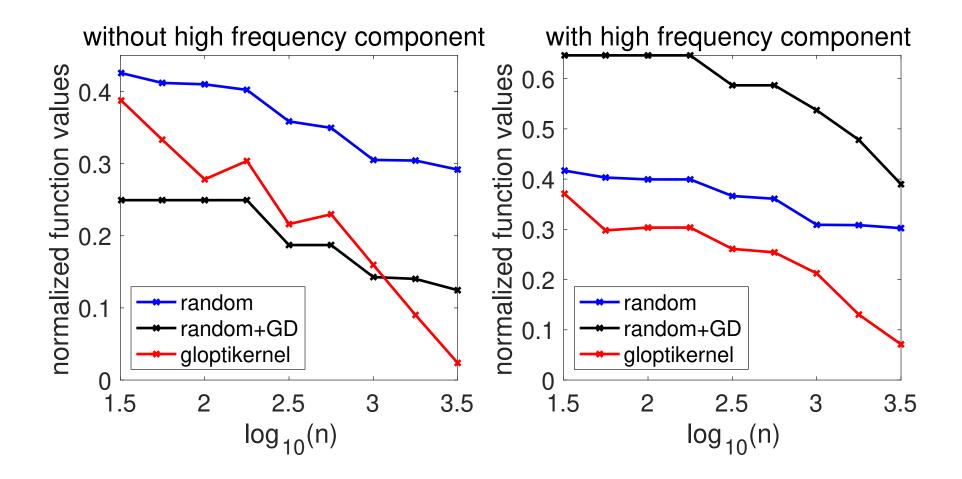


# Illustration



### Illustration

• Minimization of eight-dimensional function

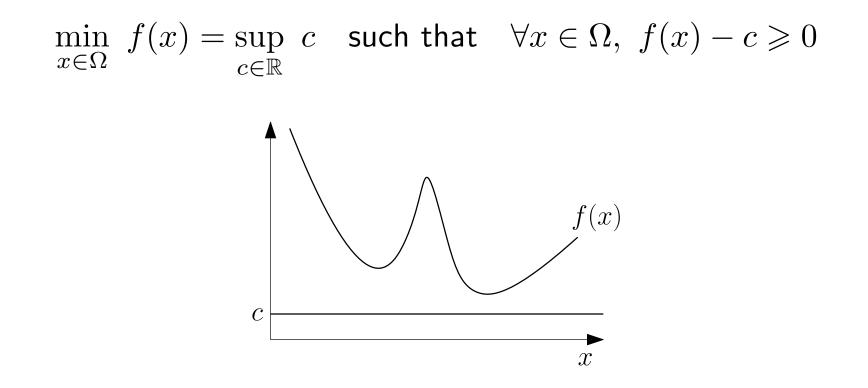


## **Extensions**

- Duality
- Extionsions
- Conclusion

# **Duality**

#### • Primal problem



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#### • Primal problem

$$\min_{x\in\Omega} f(x) = \sup_{c\in\mathbb{R}} c \quad \text{such that} \quad \forall x\in\Omega, \ f(x) - c \ge 0$$

### • Dual problem on probability measures

$$\inf_{\mu \in \mathbb{R}^{\Omega}} \int_{\Omega} \mu(x) f(x) dx \quad \text{such that} \quad \int_{\Omega} \mu(x) dx = 1, \ \forall x \in \Omega, \ \mu(x) \ge 0$$

# **Duality with sums-of-squares**

#### • Primal problem

 $\min_{x\in\Omega} f(x) = \sup_{c\in\mathbb{R},\ A\succcurlyeq 0} c \quad \text{such that} \quad \forall x\in\Omega,\ f(x) - c = \langle \phi(x), A\phi(x) \rangle$ 

• Dual problem on signed measures

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- Extension of results on polynomials (Lasserre, 2020)

## **Extension** - I

#### • Generic constrained optimization problem

$$\inf_{\theta \in \Theta} F(\theta) \quad \text{such that} \quad \forall x \in \Omega, \ g(\theta, x) \ge 0$$

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- Requires penalization by tr(A) and subsampling
- Need representation as sums-of-squares to benefit from smoothness
- Can be done in the primal or the dual

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- Requires penalization by  $\mathrm{tr}(A)$  and subsampling
- Need representation as sums-of-squares to benefit from smoothness
- Can be done in the primal or the dual
- Application to optimal transport (Vacher, Muzellec, Bach, Rudi, Vialard, 2021)

## **Extension** - II

#### • Constrained optimization problem

$$\inf_{x\in \mathbb{R}^d} f(x) \quad \text{such that} \quad \forall x\in \Omega, \ g(x) \geqslant 0$$

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#### • Sums-of-squares reformulation

$$\sup_{c \in \mathbb{R}, A \succcurlyeq 0, B \succcurlyeq 0} c$$

such that  $\forall x \in \Omega, \ f(x) = c + \langle \phi(x), A\phi(x) \rangle + g(x) \langle \phi(x), B\phi(x) \rangle$ 

- Extension of results on polynomials (Lasserre, 2001)

# Conclusion

### • Global optimization through kernel approximations

- Joint optimization and approximation
- infinite-dimensional sums-of-squares representation
- Controlled subsampling with guarantees

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### • Further extensions

- Efficient algorithms below  $O(n^3)$  complexity
- Other infinite-dimensional convex optimization problems

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- Other infinite-dimensional convex optimization problems
- See arxiv.org/abs/2012.11978 and francisbach.com/

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