

Brainstorming day on polynomial optimization:

A hierarchy of spectral relaxations for polynomial optimization

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MAC Group

joint work with my supervisors:

Jean-Bernard Lasserre and **Victor Magron**



September 10, 2020

Scalability of moment relaxations for POP

Consider a general POP:

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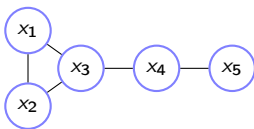
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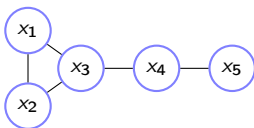
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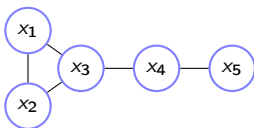
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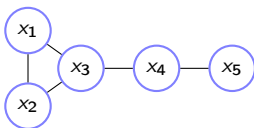
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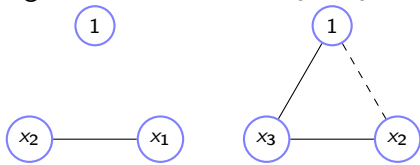
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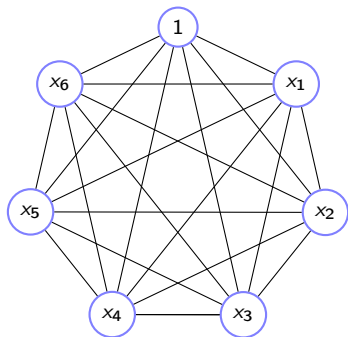
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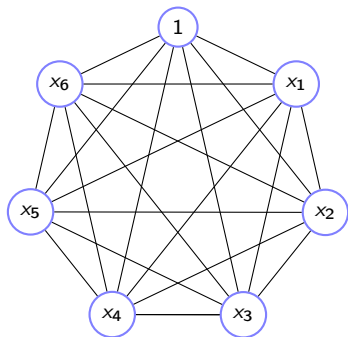
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⇒ use other information: constant trace property.

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If *constant trace property* (CTP), i.e.,

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e.g., SDP MAXCUT \Rightarrow can be solved very efficiently thanks to:

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2. SketchyCGAL=Sketchy+CG+AL: [Yurtsever et al., 2019]

- ▶ the augmented Lagrangian (AL)
- ▶ conditional gradient method (CG)
- ▶ randomized sketch (Sketchy)

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- ▶ Convert moment relaxations with CTP to spectral relaxations.
- ▶ Build an algorithm called **SpectralPOP** to solve spectral relaxations and extract the optimal solutions of POP.
- ▶ Display several numerical results of **SpectralPOP** for a sample of random dense POPs.

Moment-SOS hierarchy

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Consider an equality constrained POP on the unit sphere:



$$f^* := \inf\{f(\mathbf{x}) : h_j(\mathbf{x}) = 0, j \in [l]\} \quad \text{with} \quad h_1 = 1 - \|\mathbf{x}\|_2^2. \quad (5)$$

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$$\tau_k = \inf_{\mathbf{y}} \{ L_{\mathbf{y}}(f) : \mathbf{M}_k(\mathbf{y}) \succeq 0, \mathbf{M}_{k-w_j}(h_j \mathbf{y}) = 0, y_0 = 1 \}. \quad (6)$$

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Moreover $\tau_k = f^*$ under some assumption (see [Lasserre, 2015]).

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Obviously, the psd matrix of this form has trace 3.

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Since $\mathbf{P}_k \succ 0$, (6) is equivalent to SDP:

$$\begin{aligned} \tau_k = \inf_{\mathbf{y}} \quad & L_{\mathbf{y}}(f) \\ \text{s.t.} \quad & 2^{-k} \mathbf{P}_k \mathbf{M}_k(\mathbf{y}) \mathbf{P}_k \succeq 0, \\ & \mathbf{M}_{k-w_j}(h_j \mathbf{y}) = 0, y_0 = 1, \end{aligned} \quad (9)$$

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where $\mathcal{A}_k : \mathcal{S}_k \rightarrow \mathbb{R}^{m_k}$ is a linear operator of the form

$$\mathcal{A}_k \mathbf{X} = \left[\text{trace}(\mathbf{A}_k^{(1)} \mathbf{X}), \dots, \text{trace}(\mathbf{A}_k^{(m_k)} \mathbf{X}) \right].$$

Spectral relaxations

Following [Helmberg and Rendl, 2000], we obtain:

$$-\tau_k = \inf_{\mathbf{z}} \lambda_{\max}(\mathbf{C}_k - \mathcal{A}_k^{\top} \mathbf{z}) + \mathbf{b}_k^{\top} \mathbf{z}. \quad (11)$$

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Spectral relaxations

Following [Helmberg and Rendl, 2000], we obtain:

$$- \tau_k = \inf_{\mathbf{z}} \lambda_{\max}(\mathbf{C}_k - \mathcal{A}_k^\top \mathbf{z}) + \mathbf{b}_k^\top \mathbf{z}. \quad (11)$$

Here:

$$\mathcal{A}_k^\top \mathbf{z} = \sum_{i \in [m_k]} z_i \mathbf{A}_k^{(i)}.$$

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$m_k \geq 200 \times s_k$ when $n \geq 30$ and $k \geq 2$

\Rightarrow not suitable for spectral bundle method [Helmberg et al., 2014];

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5. Extract \mathbf{x}^* from \mathbf{y}^* ▷ [Henrion and Lasserre, 2005].

Spectral relaxations for quadratic forms ($k = 1$)

matrix size $s_k = n + 1$ 2 equality constraints

gap: the relative optimality gap w.r.t. SumOfSquares

LMBM: Limited Memory Bundle Method

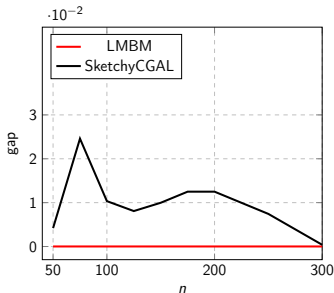
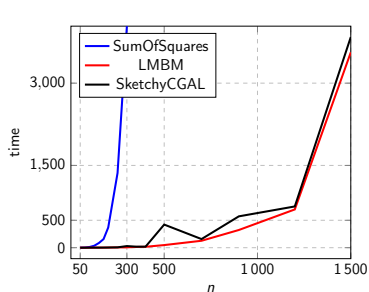
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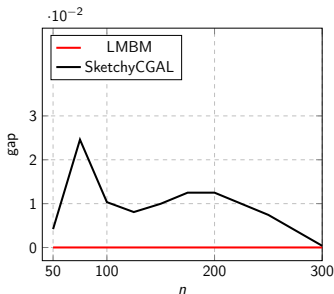
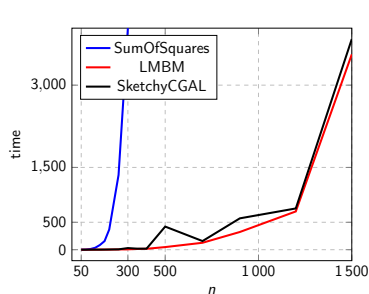
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► SpectralPOP (LMBM):

- can solve POPs with $n = 300$ variables in 5 seconds while it takes 1 hour in **SumOfSquares** (Mosek).
- provide optimal value (gap $< 10^{-5}$) and optimal solution with high accuracy for $n = 1500$ in 1 hour.

Spectral relaxations for QCQP ($k = 2$)

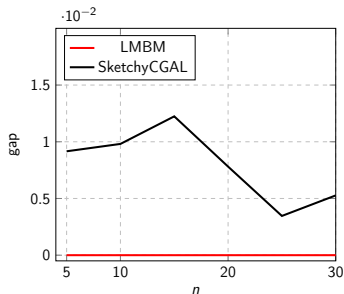
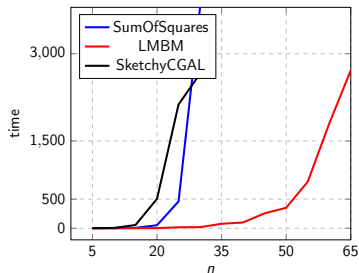
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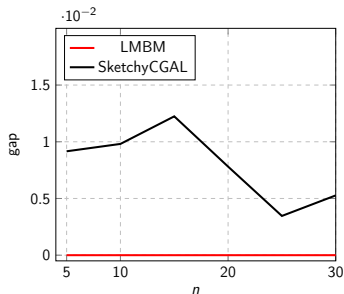
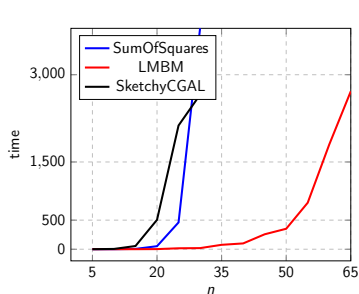
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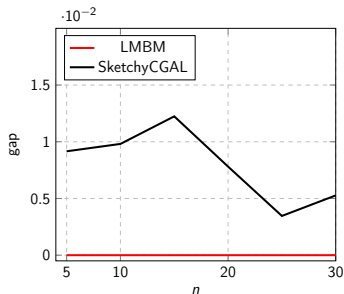
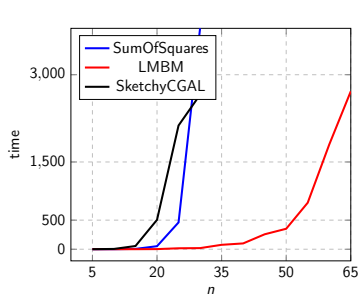


► If $n = 65$, then $s_k = 2211$ and $m_k = 1618453$.

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- ▶ If $n = 65$, then $s_k = 2211$ and $m_k = 1618453$.
- ▶ **SpectralPOP (LMBM):**
 - ▶ provide the optimal value with high accuracy (gap $< 10^{-5}$)
 - ▶ up to more than 150 times faster than **SumOfSquares** (Mosek) when $n \geq 25$.

Summary and Future works

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- ▶ Comparison with **SumOfSquares** (Mosek) and **SketchyCGAL** on the last tested sample of POPs:

SpectralPOP (LMBM) is *cheaper, faster, but maintains the same accuracy* as **SumOfSquares**.

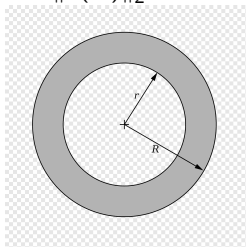
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SpectralPOP (LMBM) is *cheaper, faster, but maintains the same accuracy* as **SumOfSquares**.

- ▶ Future works:

- ▶ exploiting CTP for POPs with annulus constraints:
$$r \leq \|\mathbf{x}(T)\|_2^2 \leq R$$



- ▶ exploiting CTP for sparse POPs

End

Thank you for your attention!

SpectralPOP: <https://github.com/maihoanganh/SpectralPOP>

Homepage: <https://sites.google.com/view/hoanganhmai>



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