

Discussion on S-variable based static-output feedback design heuristics

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Find a gain $K \in \mathbb{R}^{m \times p}$ such that $\dot{x} = (A + BK C)x$ is stable

- ➔ Without having an initial guess
 \neq Knowing $A + BK_o C$ stable, find a better gain (with some criterion)
- ➔ Without having indications of a range of admissible values
 no possibility to test on a grid
- ➔ Robust w.r.t. uncertainties in A , B and C
 eg. matrices in a polytope
- ➔ Structured : eg. K diagonal (decentralized control)

Assume a linear plant $\dot{x} = A(K_1, \dots, K_{\bar{k}})x$ rational in the design parameters $K_{k=1 \dots \bar{k}}$

- ➔ Linear plant : First step before considering nonlinear dynamics
- ➔ K_k : gains of a dynamic control, filter param., decentralized control ... or plant parameters
- ➔ Rational in the parameters : or in 'gains' with one-to-one NL maps to true design parameters

It can always be reformulated (Linear Fractional Transformation) as a feedback control loop

$$\begin{cases} \dot{x} = Ax + Bu \\ y = Cx + Du \end{cases} \quad u = Ky = \begin{bmatrix} I_{r_1} \otimes K_1 & & 0 \\ & \ddots & 0 \\ 0 & & I_{r_{\bar{k}}} \otimes K_{\bar{k}} \end{bmatrix} y$$

where K is structured (block-diagonal), parameters may be repeated: $I_{r_k} \otimes K_k$.

- ➔ with $y := y - Du$ one may consider $D = 0$
- ➔ All these design problems look 'simple': find a (structured) K s.t. $A + BKC$ is Hurwitz

Find a (structured) K s.t. $A + BKC$ is Hurwitz \Leftrightarrow minimization of a non-linear, non-smooth function

- ➔ [Apkarian, Noll 2003] optimize an H_∞ gain (and more) - hinfstruct - Clark's sub-gradient
- ➔ [Overton et al. 2006] optimize an H_∞, H_2 gains, spectral abscissa - Hifoo - gradient sampling
- ➔ [Peretz 2013] optimize the spectral abscissa - randomized approximation

▲ ▲ Very efficient in practice

▼ Randomized flavour (random initial conditions, randomization in the algorithm)

▼ No extensions to robust design

▲ Lyapunov based approaches with matrix inequalities may handle robustness

$$P \succ 0 \quad , \quad (A + BKC)^* P + P(A + BKC) = \{P(A + BKC)\}^{\mathcal{H}} \prec 0$$

$$Q \succ 0 \quad , \quad \{(A + BKC)Q\}^{\mathcal{H}} \prec 0$$

▼ Bilinear matrix inequalities (not convex)

Find a (structured) K s.t. $A + BKC$ is Hurwitz is sometimes convex (up to a transformation)

- ➔ $B = C = I$ State-Injection : $\{PA + L\}^{\mathcal{H}} \prec 0$ gives $K = P^{-1}L$
- ➔ $B = I$ Output-Injection : $\{PA + LC\}^{\mathcal{H}} \prec 0$ gives $K = P^{-1}L$
- ➔ $C = I$ State-Feedback : $\{AQ + BF\}^{\mathcal{H}} \prec 0$ gives $K = FQ^{-1}$
- ➔ Almost commutative on the input : $\{PA + BLC\}^{\mathcal{H}} \prec 0$, $PB = B\hat{P}$ gives $K = \hat{P}^{-1}L$
- ➔ Almost commutative on the output : $\{AQ + BFC\}^{\mathcal{H}} \prec 0$, $CQ = \hat{Q}C$ gives $K = F\hat{Q}^{-1}$

▼ Applies only to special cases

▲ Robustness can be considered

(eg. test on all vertices with common decision variables proves stability of the polytope)

▼ Structured K cannot be considered

unless one considers structured P and Q (very conservative)

Simple P - K -iterative algorithm

- Initialization Choose a positive definite P
- K-iteration For fixed P find $K = \arg \min \alpha$ under $\{P(A + BK C)\}^{\mathcal{H}} \prec \alpha I$
- P-iteration For fixed K find $P = \arg \min \alpha$ under $\{P(A + BKC)\}^{\mathcal{H}} \prec \alpha I$
- Stop Repeat until $\alpha < 0$ (success) or varies too slowly (failure)

- ▲ Easy to implement
- ▲ Strictly decreasing sequence of α
- ▼ Very sensitive to initialization
- ▼ Little progress after very few steps
- ▼ Not effective in practice

$$\exists S : M \prec \{X_1^* S X_2\}^{\mathcal{H}} \quad \Leftrightarrow \quad \begin{cases} N_{X_1}^* M N_{X_1} \prec 0 \\ N_{X_2}^* M N_{X_2} \prec 0 \end{cases} : X N_X = 0, \text{ Rank}(N_X) = \dim(\text{Ker}(X))$$

Applied to the SOF problem [Scherer, Iwasaki...] it gives

$$\exists K : \begin{cases} \{P(A + BK C)\}^{\mathcal{H}} \prec 0 \\ \{(A + BK C)Q\}^{\mathcal{H}} \prec 0 \\ PQ = I \end{cases} \quad \Leftrightarrow \quad \begin{cases} N_C^* \{PA\}^{\mathcal{H}} N_C \prec 0 \\ N_{B^*}^* \{AQ\}^{\mathcal{H}} N_{B^*} \prec 0 \\ PQ = I \end{cases}$$

- ▲ May be converted to pure LMIs (of small dimensions) ... ▼ with a rank constraint
- ▲ Many dedicated iterative algorithms
- ▼ Sensitive to initial guesses, not very effective in practice
- ▼ Cannot take into account structured K

$$\exists S : M \prec \{X_1^* S X_2\}^{\mathcal{H}} \quad \Leftrightarrow \quad \begin{cases} N_{X_1}^* M N_{X_1} \prec 0 \\ N_{X_2}^* M N_{X_2} \prec 0 \end{cases} : X N_X = 0 \quad \text{Rank}(N_X) = \dim(\text{Ker}(X))$$

Assume P proves stability for both an output-feedback gain K_{of} and a state-feedback gain K_{sf}
 (always true with $K_{sf} = K_{of} C$)

$$P(A + B K_{of} C)^{\mathcal{H}} = \begin{bmatrix} I \\ K_{of} C \end{bmatrix}^* \begin{bmatrix} \{PA\}^{\mathcal{H}} & PB \\ B^* P & 0 \end{bmatrix} \begin{bmatrix} I \\ K_{of} C \end{bmatrix} \prec 0$$

$$P(A + B K_{sf})^{\mathcal{H}} = \begin{bmatrix} I \\ K_{sf} \end{bmatrix}^* \begin{bmatrix} \{PA\}^{\mathcal{H}} & PB \\ B^* P & 0 \end{bmatrix} \begin{bmatrix} I \\ K_{sf} \end{bmatrix} \prec 0$$

Equivalent to the existence of S such that

$$\begin{bmatrix} \{PA\}^{\mathcal{H}} & PB \\ B^* P & 0 \end{bmatrix} \prec \left\{ \begin{bmatrix} K_{sf}^* \\ -I \end{bmatrix} S \begin{bmatrix} K_{of} C & -I \end{bmatrix} \right\}^{\mathcal{H}}$$

$$\exists P \succ 0, S, K_{sf}, K_{of} : \begin{bmatrix} \{PA\}^{\mathcal{H}} & PB \\ B^*P & 0 \end{bmatrix} \prec \left\{ \begin{bmatrix} K_{sf}^* \\ -I \end{bmatrix} S \begin{bmatrix} K_{of}C & -I \end{bmatrix} \right\}^{\mathcal{H}}$$

- ▼ Still not LMI, matrix inequalities of larger size, more decision variables
- ▲ More degrees of freedom, P and K_{of} separated one for the other
- ▲ Simple to code K_{sf} - K_{of} -iterative algorithm

$$K_{sf}\text{-iteration} \quad K_{of} = S^{-1}L_{of} \quad \arg \min \alpha : \begin{bmatrix} \{PA\}^{\mathcal{H}} - \alpha I & PB \\ B^*P & 0 \end{bmatrix} \prec \left\{ \begin{bmatrix} K_{sf}^* \\ -I \end{bmatrix} \begin{bmatrix} L_{of}C & -S \end{bmatrix} \right\}^{\mathcal{H}}$$

$$K_{of}\text{-iteration} \quad K_{sf} = S^{-1}L_{sf} \quad \arg \min \alpha : \begin{bmatrix} \{PA\}^{\mathcal{H}} - \alpha I & PB \\ B^*P & 0 \end{bmatrix} \prec \left\{ \begin{bmatrix} L_{sf}^* \\ -S \end{bmatrix} \begin{bmatrix} K_{of}C & -I \end{bmatrix} \right\}^{\mathcal{H}}$$

- ▲ Smart initial guess of K_{sf} (finding K_{sf} is a convex problem)
- ▲ Strictly decreasing sequence of α
- ▲ Much more efficient than the P - K -iterative algorithm (P is free at each step)
 Implicitly the algorithm searches for $K_{sf} \rightarrow K_{of}C$.

Variant of the same result based on output-injection gain K_{oi}

$$\exists Q \succ 0, S, K_{oi}, K_{of} \quad : \quad \begin{bmatrix} \{AQ\}^{\mathcal{H}} & QC^* \\ CQ & 0 \end{bmatrix} \prec \left\{ \begin{bmatrix} BK_{of} \\ -I \end{bmatrix} S \begin{bmatrix} K_{oi}^* & -I \end{bmatrix} \right\}^{\mathcal{H}}$$

- ▲ Simple to code K_{oi} - K_{of} -iterative algorithm
- ▲ Smart initial guess of K_{oi} (finding K_{oi} is a convex problem)
- ▲ Strictly decreasing sequence of α
- ▲ Much more efficient than the P - K -iterative algorithm (P is free at each step)
 Implicitly the algorithm searches for $K_{oi} \rightarrow BK_{of}$.
- ▲ Robustness can be dealt with easily (eg. solve the constraints for all vertices of a polytope)
 - ▼ yet conservative (common Lyapunov certificate P or Q for all uncertainties)
- ▲ Structured SOF : achievable with constraints on S , not on P of Q
 - ▼ yet conservative

Variant [Peres et al. 2020] assuming two Lyapunov certificates P_1 and P_2 for $A + BK C$

$$\exists P_1 \succ 0, P_2 \succ 0, S = S^*, K \quad : \quad \begin{bmatrix} 0 & (A + BK C)^* \\ A + BK C & 0 \end{bmatrix} \prec \left\{ \begin{bmatrix} P_2 \\ -I \end{bmatrix} S \begin{bmatrix} P_1 & -I \end{bmatrix} \right\}^{\mathcal{H}}$$

- ▲ Matrix inequalities of larger size
- ▲ Simple to code P_1 - P_2 -iterative algorithm
- ▼ No smart initial guess of P
- ▲ Robustness can be dealt with easily (eg. solve the constraints for all vertices of a polytope)
- ▲ No difficulty to include structure constraints on K
- ▼ Seems less efficient than the K_{sf} - K_{of} -iterative algorithm

$$\begin{bmatrix} 0 & 0 & Q \\ 0 & 0 & 0 \\ Q & 0 & 0 \end{bmatrix} \prec \left\{ \begin{bmatrix} -I \\ LC + M \\ A \end{bmatrix} S_1 + \begin{bmatrix} 0 \\ -S_2 \\ BF \end{bmatrix} [0 \quad -I \quad H^*] \right\}^{\mathcal{H}}$$

- ➔ If $L = 0$ then $K_{si} = HM$ is a stabilizing state-injection gain ($A + K_{si}$ is stable)
- ➔ If $L = 0$ then $K_{sf} = FS_2^{-1}M$ is a stabilizing state-feedback gain ($A + BK_{sf}$ is stable)
- ➔ If $M = 0$ then $K_{oi} = HL$ is a stabilizing output-injection gain ($A + K_{oi}C$ is stable)
- ➔ If $M = 0$ then $K_{of} = FS_2^{-1}L$ is a stabilizing output-feedback gain ($A + BK_{of}C$ is stable)

▲ Stabilizing state-injection property : Good for initialization

▲ All matrices Q , A , B and C decoupled:

OK for robustness with parameter-dependent Lyapunov certificates

▲ Results are new (and efficient) even for robust K_{sf} and K_{oi} design

▲ No need to structure the Lyapunov certificate for structured SOF

▼ Algorithm is less trivial than the previous ones (contains a line search)

➔ (S_1, K_{sf}) - K_{si} -iterative algorithm shows to be efficient with low number of iterations

➔ During the algorithm $K_{si} \rightarrow K_{oi}C$ and $K_{sf} \rightarrow K_{of}C$

- ➔ Several results for SOF using the S-variable approach
 - ▲ Rather efficient to deal with robust structured SOF
 - ▼ No guarantee of success
 - ▲ ▼ Results with more or less intuitive initialization

- ➔ Latest result - In Honor of Roberto Tempo
 - ▲ Promising numerical experiments
 - ▲ Elegant combination of the SI, OI, SF, OF problems
 - ▲ Variations of the algorithms for
 - Deterministic robust design
 - Probabilistic robust design
 - With comparisons of the two

"Robust static output feedback design with deterministic and probabilistic certificates",
D. Arzelier, F. Dabbene, S. Formentin, D. Peaucelle, L. Zaccarian
Birkhäuser Mathematics - Springer Nature Chapter for "Uncertainty in Networked Systems"