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Robust anti-windup 0000000000 Non-common certificate

References

Randomized robust static anti-windup

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> Séminaire MOSAR Onera, Toulouse, November 28, 2014



- Saturation: an abrupt nonlinearity:
 - Small signals: $sat(u) = u \Rightarrow$ no effect
 - Large signals: sat(u) bounded \Rightarrow severe effect

• Signal size (
$$\mathcal{L}_2$$
 norm): $\|z\|_2 := \left(\int_0^\infty |z(t)|^2 dt\right)^{rac{1}{2}}$

- $z \in \mathcal{L}_2$ (square integrable) if $\|z\|_2 < \infty$
- Closed-loop performance measures:
 - Finite \mathcal{L}_2 gain (linear \mathcal{H}_∞ norm): $\overline{\gamma}_{wz} \in \mathbb{R}_{\geq 0}$:

 $\|z\|_2 \leq \overline{\gamma}_{wz} \|w\|_2$ for all $w \in \mathcal{L}_2$

• Nonlinear \mathcal{L}_2 gain: a function $s \mapsto \gamma_{wz}(s)$: Megretski [1996] $\|z\|_2 \le \gamma_{wz}(s) \|w\|_2$ for all w satisfying $\|w\|_2 \le s$



Non-common certificates

References

Example demonstrates relevance of nonlinear gains

Controller K cancels the plant dynamics and stabilizes (before saturation)

$$\mathcal{P}: \dot{z} = az + \operatorname{sat}(u) + w$$
$$\mathcal{K}: u = -az - 10z$$



Three representative cases Sontag [1984], Lasserre [1992]



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Robust anti-windup

Non-common certificates

References

Optimal nominal static linear anti-windup design (LMI)



- Given \mathcal{P} linear, \mathcal{C} linear, design only
 - linear anti-windup gain $D_{aw} = \begin{bmatrix} D_{aw,1} \\ D_{aw,2} \end{bmatrix}$
- Performance objective:
 given s*, minimize γ_{dz}(s*)
- Linear controller \mathcal{K} equations $\dot{x}_c = Ax_c + By + D_{aw,1}(u - \operatorname{sat}(u))$ $y_c = Cx_c + Dy + D_{aw,2}(u - \operatorname{sat}(u))$
- LMI-based design Mulder et al. [2001], Gomes da Silva Jr and Tarbouriech [2005], Hu et al. [2008]
- Preserve of small signal response (*D_{aw}* multiplies dz(*u*) = *u* sat(*u*))
 Asymptotically recover large signal response (global not always possible)
- Robust designs follow a deterministic worst case paradigm, imposing strong convexity conditions Turner et al. [2007], Grimm et al. [2004]
- This talk: randomized analysis and synthesis of robust static anti-windup

• Quadratic functions (LMIs Boyd et al. [1994])

 $V_1(x) = x^T P x$

- Max of quadratics (BMIs) $V_2(x) = \max_{j \in \{1,...,J\}} x^T P_j x$
- Convex Hull of quadratics (BMIs) $V_3(x) = \min_{\gamma_j \ge 0: \sum_j \gamma_j = 1} x^T \left(\sum_j \gamma_j Q_j \right)^{-1} x$
- **Piecewise** quadratic (LMI-BMI) $V_4(x) = \begin{bmatrix} x \\ dz(u(x)) \end{bmatrix}^T \overline{P} \begin{bmatrix} x \\ dz(u(x)) \end{bmatrix}$
- Piecewise **Polynomial** (LMI-BMI) $V_{5}(x) = \begin{bmatrix} x \\ dz(u(x)) \end{bmatrix}^{\{m\}T} \hat{P} \begin{bmatrix} x \\ dz(u(x)) \end{bmatrix}^{\{m\}T}$

$$\dot{V}+rac{1}{\gamma_{dz}(s)}|z|^2-\gamma_{dz}(s)|w|^2<0$$



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$$V_1(x) = x^T P x$$

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$$\dot{\lambda} + rac{1}{\gamma_{dz}(s)}|z|^2 - \gamma_{dz}(s)|w|^2 < 0$$



Nominal analysis and design

Robust anti-windup

Non-common certificate

References

Compact representation of the closed-loop system



$$\mathcal{H}: \begin{cases} \dot{x_{cl}} = A_{cl}x_{cl} + B_{cl,d}(u - \operatorname{sat}(u)) + B_{cl,v}v + B_{cl,w}w \\ u = C_{cl,u}x_p + D_{cl,ud}(u - \operatorname{sat}(u)) + D_{cl,uv}v + D_{cl,uw}w \\ z = C_{cl,z}x_p + D_{cl,zd}\underbrace{(u - \operatorname{sat}(u))}_{\operatorname{dz}(u)} + D_{cl,zv}v + D_{cl,zw}w, \end{cases}$$

Proposition: Given the NOMINAL system and s > 0, if the LMI problem

$$\begin{split} \hat{\gamma}^{2}(s) &= \min_{\{\gamma^{2}, Q, Y, U\}} \gamma^{2} \text{ subject to } Q = Q^{T} > 0, \ U > 0 \text{ diagonal}, \\ \mathrm{He} \begin{bmatrix} A_{cl}Q & B_{cl,d}U + B_{cl,v}D_{aw}U + Y^{T} & B_{cl,w} & 0\\ C_{cl,u}Q & D_{cl,ud}U + D_{cl,uv}D_{aw}U - U & D_{cl,uw} & 0\\ 0 & 0 & -I/2 & 0\\ C_{cl,z}Q & D_{cl,zd}U + D_{cl,zv}D_{aw}U & D_{cl,zw} - \frac{\gamma^{2}}{2}I \end{bmatrix} \prec 0, \ \begin{bmatrix} Q & Y_{[k]}^{T} \\ Y_{[k]} & \overline{u}_{k}^{2}/s^{2} \\ k = 1, \dots, n_{u} \end{bmatrix} \succeq 0 \end{split}$$

is feasible, then the following holds for the saturated closed-loop:

- [Stab] the origin is locally exponentially stable with region of attraction containing the set *E*(*Q*, *s*) := {*x* : *x*^T*Q*⁻¹*x* ≤ *s*²};
- [Reach] the reachable set from x(0) = 0 with ||w||₂ ≤ s is contained in *E*(Q, s);
- [L₂Perf] for each w such that ||w||₂ ≤ s, the zero state solution satisfies the L₂ gain bound:

$$\|z\|_2 \leq \hat{\gamma}(s) \|w\|_2$$

Saturation and anti-windup OCO COOCOCOO Quadratic analysis conditions easily lead to synthesis Mulder et al. [2001], Gomes da Silva Jr and Tarbouriech [2005], Hu et al. [2008]

Proposition: Given the NOMINAL system and s > 0. If the LMI problem

$$\hat{\gamma}^{2}(s) = \min_{\{\gamma^{2}, Q, Y, U\}} \gamma^{2} \text{ subject to } Q = Q^{T} > 0, \ U > 0 \text{ diagonal},$$

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is feasible, then, selecting the static AW gain as

 $D_{aw} = XU^{-1}$

- [Stab] the origin is locally exponentially stable with region of attraction containing the set *E(Q, s)* := {*x* : *x^TQ⁻¹x* ≤ *s*²};
- [Reach] the reachable set from x(0) = 0 with ||w||₂ ≤ s is contained in E(Q, s);
- [L₂Perf] for each w such that ||w||₂ ≤ s, the zero state solution satisfies the L₂ gain bound:

 $\|z\|_2 \leq \hat{\gamma}(s) \|w\|_2$



$$\mathcal{H}(q): \begin{cases} \dot{x_{cl}} = A_{cl}(q)x_{cl} + B_{cl,d}(q)(u - \operatorname{sat}(u)) + B_{cl,v}(q)v + B_{cl,w}(q)w \\ u = C_{cl,u}(q)x_p + D_{cl,ud}(q)(u - \operatorname{sat}(u)) + D_{cl,uv}(q)v + D_{cl,uw}(q)w \\ z = C_{cl,z}(q)x_p + D_{cl,zd}(q)\underbrace{(u - \operatorname{sat}(u))}_{\operatorname{dz}(u)} + D_{cl,zv}(q)v + D_{cl,zw}(q)w, \end{cases}$$



To solve the robust synthesis problem, may look for θ = {γ², Q, Y, U, X}
 s.t. the LMI holds for all q ∈ Q

Given a scalar s > 0, if the **nonconvex** optimization problem is feasible

 $\hat{\gamma}^2(s) = \min_{ heta} \gamma^2$ subject to $Q = Q^T > 0, \ U > 0$ diagonal,

$$\operatorname{He} \begin{bmatrix} A_{cl}(q)Q & B_{cl,d}(q)U + B_{cl,v}(q)X + Y^{T} & B_{cl,w}(q) & 0\\ C_{cl,u}(q)Q & D_{cl,ud}(q)U + D_{cl,uv}(q)X - U & D_{cl,uw}(q) & 0\\ 0 & 0 & -I/2 & 0\\ C_{cl,z}(q)Q & D_{cl,zd}(q)U + D_{cl,zv}(q)X & D_{cl,zw}(q) - \frac{\gamma^{2}}{2}I \end{bmatrix} \prec 0, \quad \forall q \in \mathbb{Q}$$

$$\begin{bmatrix} Q & Y_{[k]}^{T} \\ Y_{[k]} & \overline{u}_{k}^{2}/s^{2} \end{bmatrix} \succeq 0, \quad k = 1, \dots, n_{u}$$

then, selecting the static AW gain as

$$D_{aw} = XU^{-1}$$

[WP], [LocStab], and [\mathcal{L}_2 Perf] are robustly guaranteed

This construction is hard due to general dependence on *q*.
 ⇒ Can use scenario (or sequential) randomized approach



To solve the robust synthesis problem, may look for θ = {γ², Q, Y, U, X}
 s.t. the LMI holds for all q ∈ Q

Given a scalar s > 0, if the **nonconvex** optimization problem is feasible $\hat{\gamma}^2(s) = \min_{\theta} c^{\top} \theta$ subject to $f_s(\theta, q) \le 0, \forall q \in \mathbb{Q}$ where $\theta \in \mathbb{R}^{n_{\theta}}$ are the design variables; $q \in \mathbb{Q}$ are the uncertain parameters;

• $f_s(\theta, q) \leq 0$ are the problem constraints

then, selecting the static AW gain as

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[WP], [LocStab], and [\mathcal{L}_2 Perf] are robustly guaranteed

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| Saturation and anti-windup 000 | Nominal analysis and design | | Robust ant 000€00 | ti-win |
|-----------------------------------|-----------------------------|---|----------------------|--------|
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Non-common certificates

References

Robust and chance constrained optimization

Problem (Robust optimization - RO)

Given an objective vector $c \in \mathbb{R}^{n_{\theta}}$, solve

 $\min_{\theta} c^{\top} \theta \qquad subject \ to$ $f(\theta, q) < 0, \text{ for all } q \in \mathbb{O}$

Problem (Chance constrained optimization - CC)

Let a distribution over \mathbb{Q} be given, and let $\varepsilon \in (0, 1)$ be a (small) probability level. Given an objective vector $c \in \mathbb{R}^{n_{\theta}}$, solve

$$\min_{\theta} c^{\top} \theta \qquad subject \ to$$
$$\underbrace{\operatorname{Prob}\{q \in \mathbb{Q} : f(\theta, q) > 0\}}_{\operatorname{Viol}(\theta)} \leq \varepsilon$$

| Saturation and anti-windup | Nominal analysis and design | Robust anti-windup | Non-common certificates | References |
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Calafiore et al. [2011], Tempo et al. [2013], Petersen and Tempo [2014]

Both problems are very hard in general

- Robust optimization is hard whenever the uncertainty enters in a nonlinear way
- Chance-constrained optimization is a even more difficult non-convex problem (it involves hard integral evaluations)
- Proposed solution approach: Randomized algorithms
 - A Randomized Algorithm is an algorithm that makes random choices during its execution to produce a result



• Randomized algorithms entail a (pre determined) probability of failure

Nominal analysis and desig

Robust anti-windup

Non-common certificates

References

Scenario approach amounts to a simple extraction Calafiore and Campi [2006]

- Scenario techniques provide a simple and theoretically sound way to approximately solve the two problems RO and CC
- The idea is to replace these hard optimization problems with the following sampled counterpart (random convex program)

Problem (Scenario optimization)

Extract N i.i.d. samples (scenarios) $q^{(1)}, \ldots, q^{(N)}$, and solve

$$\min_{\theta} c^{\top} \theta \quad subject \ to \\ f(\theta, q^{(i)}) \leq 0, \quad i = 1, \dots, N$$

• The scenario problem is a **standard convex optimization problem** with a finite number of constraints

Nominal analysis and desig

Robust anti-windup

Non-common certificates

References

Probability of violation is bounded by number of samples

Assumption (Basic assumptions)

 $f(\theta, q)$ is continuous and convex in θ for any fixed $q \in \mathbb{Q}$. For any multisample extraction $\mathbf{q} = \{q^{(1)}, \ldots, q^{(M_k)}\} \in \mathbb{Q}$, the scenario problem is feasible and attains a unique optimal solution

Theorem (violation of scenario solutions Campi and Garatti [2008])

Let $\varepsilon \in (0, 1)$ be a given probability level and let $N \ge n_{\theta}$. Under **convexity**, **uniqueness** and **feasibility** assumptions, the scenario solution θ_{sc} satisfies

$$\mathsf{Pr}\big\{\mathsf{Viol}(\theta_{sc}) > \varepsilon\big\} \le \mathsf{B}(N,\varepsilon, \mathbf{n}_{\theta})$$

where

$$\mathsf{B}(\mathsf{N},\varepsilon,\mathsf{n}_{\theta})=\sum_{k=0}^{\mathsf{n}_{\theta}-1}\binom{\mathsf{N}}{k}\varepsilon^{k}(1-\varepsilon)^{\mathsf{N}-k}.$$

Nominal analysis and desig

Robust anti-windup

Non-common certificates

References

Robust static anti-windup synthesis based on scenario

Theorem (Robust static AW synthesis using scenario Formentin et al. [2013])

Fix a positive value $s \ge \|w\|_2$, $\varepsilon \in (0, 1)$, $\beta \in (0, 1)$, and select N satisfying

 $\mathsf{B}(N,\varepsilon, \mathbf{n}_{\theta}) \leq \beta,$

with $n_{\theta} = 1 + n(n+1)/2 + nn_u + n_u + n_u (n_u + n_c)$

Extract *N* samples of the uncertain matrices according to the probability distribution **Solve**

$$\begin{split} & \gamma_{sc}^{2}(s) = \min_{\left\{\gamma^{2}, Q, Y, U, X\right\}} \gamma^{2}, \quad subject \ to \ Q = Q^{T} > 0, \ U > 0 \ diagonal, \\ & \operatorname{He} \begin{bmatrix} A_{cl}^{(i)} Q & B_{cl,d}^{(i)} U + B_{cl,v}^{(i)} X + Y^{T} & B_{cl,w}^{(i)} & 0 \\ C_{cl,u}^{(i)} Q & D_{cl,ud}^{(i)} U + D_{cl,uv}^{(i)} X - U & D_{cl,uw}^{(i)} & 0 \\ 0 & 0 & -I/2 & 0 \\ C_{cl,z}^{(i)} Q & D_{cl,zd}^{(i)} U + D_{cl,zv}^{(i)} X & D_{cl,zw}^{(i)} & -\frac{\gamma^{2}}{2}I \end{bmatrix} < 0, \quad \begin{bmatrix} Q & Y_{[k]}^{T} \\ Y_{[k]} & \overline{u}_{k}^{2}/s^{2} \\ \forall k = 1, \dots, n_{u} \\ \forall i = 1, \dots, N \end{split}$$

If the above LMIs are feasible, select the static anti-windup gain $D_{aw} = XU^{-1}$

Then, for each $\|w\|_2 < s$, the zero initial state solution of the closed loop satisfies $Pr(\|z\|_2 > \gamma_{sc}(s) \|w\|_2) \le \varepsilon$, with probability no smaller than $1 - \beta$.

Analysis conditions can also be easily formulated

Nominal analysis and design

Robust anti-windup

Non-common certificates

References

Illustrative example: a double RC passive network



• Uncertain parameters with (known) Gaussian distribution

| parameter | mean | standard deviation |
|-----------------------|--------|--------------------|
| R_1 | 310 Ω | \pm 10 % |
| R_2 | 10 Ω | \pm 10 % |
| <i>C</i> ₁ | 0.01 F | \pm 10 % |
| <i>C</i> ₂ | 0.01 F | \pm 10 % |

- Input generator voltage constrained: $u(t) \in [-\bar{u}, \bar{u}] = [-1 Volt, 1 Volt]$
- Design parameters are $\varepsilon = 0.05, \beta = 10^{-6}, s = 1$ $\Rightarrow N = 1323$ for analysis and N = 1482 for synthesis



- The probabilistic robust compensator shows better performance (left curves)
- The nominal behavior slightly deteriorated (right curves)



Without anti-windup (black solid), with nominal anti-windup (blue dashed) and with robust anti-windup (red dashed-dotted)

| Saturation and anti-windup | Nominal analysis and design | Robust anti-windup | Non-common certificates | Reference |
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- The scenario approach to AW provides a new viewpoint to robust AW design *allowing us to address hard nonconvex synthesis problems*
- However, it is still very conservative, because we are looking for a common quadratic Lyapunov function x^T Q⁻¹x for all q ∈ Q that is for "common certificates" of stability and performance
- We would like to have "parameter dependent certificates" because non-common Lyapunov functions are known to lead to greatly reduced conservatism
- Indeed, a (much) less conservative solution can be obtained by looking for design variables θ = {γ², U, X} such that, for each q ∈ Q, there exist certificates ξ = {Q, Y} = {Q(q), Y(q)} satisfying the stability/performance LMIs
- This approach is new to within the randomized world. We denote it **design with certificates**



Given a scalar s > 0, if the LMI problem

 $\hat{\gamma}^{2}(s) = \min_{\left\{\gamma^{2}, U, X, Q, \gamma\right\}} \gamma \text{ subject to } U > 0 \text{ diagonal},$

$$\begin{split} Q &= Q^{T} > 0 \\ \mathrm{He} \begin{bmatrix} A_{cl}(q)Q & B_{cl,d}(q)U + B_{cl,v}(q)X + Y^{T} & B_{cl,w}(q) & 0 \\ C_{cl,u}(q)Q & D_{cl,ud}(q)U + D_{cl,uv}(q)X - U & D_{cl,uw}(q) & 0 \\ 0 & 0 & -I/2 & 0 \\ C_{cl,z}(q)Q & D_{cl,zd}(q)U + D_{cl,zv}(q)X & D_{cl,zw}(q) & -\frac{\gamma^{2}}{2}I \end{bmatrix} \prec 0, \ \forall q \in \mathbb{Q} \\ \begin{bmatrix} Q & Y_{[k]}^{T} \\ Y_{[k]} & \overline{u}_{k}^{2}/s^{2} \end{bmatrix} \succeq 0, \quad k = 1, \dots, n_{u} \end{split}$$

is feasible, then, selecting the static AW gain as

$$D_{aw} = XU^{-2}$$

all properties [Stab], [Reach], and [\mathcal{L}_2 Perf] hold robustly

 Saturation and anti-windup
 Nominal analysis and design
 Robust anti-windup
 Non-common certificates
 Reference

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 Robust static
 AW synthesis – non-common certificates
 Reference

Given a scalar s > 0, if the LMI problem

 $0 = 0^T \ge 0$

$$\hat{\gamma}^2(s) = \min_{\left\{\gamma^2, U, X
ight\}} \gamma$$
 subject to $U > 0$ diagonal,

for each $q \in \mathbb{Q}$ there exist $\{Q_q, Y_q\}$ such that

$$\begin{split} & \operatorname{He} \begin{bmatrix} A_{cl}(q)Q_q & B_{cl,d}(q)U + B_{cl,v}(q)X + Y_q^T & B_{cl,w}(q) & 0\\ C_{cl,u}(q)Q_q & D_{cl,ud}(q)U + D_{cl,uv}(q)X - U & D_{cl,uw}(q) & 0\\ 0 & 0 & -I/2 & 0\\ C_{cl,z}(q)Q_q & D_{cl,zd}(q)U + D_{cl,zv}(q)X & D_{cl,zw}(q) - \frac{\gamma^2}{2}I \end{bmatrix} \prec 0, \\ & \begin{bmatrix} Q_q & Y_{q[k]}^T \\ Y_{q[k]} & \overline{u}_k^2/s^2 \end{bmatrix} \succeq 0, \quad k = 1, \dots, n_u \end{split}$$

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| Saturation and anti-windup | Nominal analysis and design | Robust anti-windup 0000000000 | Non-common certificates | Reference |
|----------------------------|-----------------------------|----------------------------------|-------------------------|-----------|
| Robust optin | nization with c | ertificates | | |
| Problem (Ro | bust optimization wit | ch certificates Ois | shi [2006]) | |
| | min $c^T \theta$ su | bject to | (F | (wC) |

$$\theta \in \mathcal{S}(q), \text{ for all } q \in \mathbb{Q},$$

where the set S(q) is defined as

 $S(q) \doteq \{\theta \in \mathbb{R}^{n_{\theta}} \text{ such that there exists } \xi \text{ satisfying } f(\theta, \xi, q) \leq 0\}.$

- The idea of constructing certificates based on random samples was originally introduced by Oishi [2006], in the context of randomized ellipsoid method
- Essentially, one is allowed to use "parameter-dependent" certificates $\xi = \xi(q)$ (e.g., parameter-dependent Lyapunov functions)
- The scenario with certificates approach allows to find a solution *without* explicitly assuming the form of the dependence

 Saturation and anti-windup
 Nominal analysis and design
 Robust anti-windup
 Non-common certificates
 Reference

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 Reference

 Approximate RwC based on multisample extraction

Scenario with certificates: contrary to the scenario problem, now a new certificate variable ξ_i is used for each sample $q^{(i)}$, i = 1, ..., N

Problem (Scenario with certificates Formentin et al. [2014])

 $\min_{\theta,\xi_1,\ldots,\xi_N} c^{\mathsf{T}}\theta \text{ subject to:}$

Theorem (Scenario with certificates Formentin et al. [2014])

If for any multisample extraction the SwC problem is feasible and attains a unique optimal solution θ_{swc} , then, given an accuracy level $\varepsilon \in (0, 1)$, the solution θ_{swc} satisfies

 $f(\theta, \xi_i, \boldsymbol{q}^{(i)}) < 0, \ \forall i = 1, \dots, N$

 $\Pr\left\{\mathsf{Viol}(\theta_{\mathsf{swc}}) > \varepsilon\right\} \le \mathsf{B}(\mathsf{N}, \varepsilon, \mathsf{n}_{\theta})$

A sequential algorithm for SwC is also presented in Formentin et al. [2014]

(SwC)

Nominal analysis and desig

Robust anti-windup

Non-common certificates

References

Static AW synthesis based on scenario (recall)

Theorem (Robust static AW synthesis using scenario Formentin et al. [2013])

Fix a positive value s $\geq \|w\|_2, \, \varepsilon \in (0,1), \, \beta \in (0,1),$ and select N satisfying

 $\mathsf{B}(N,\varepsilon, \mathbf{n}_{\theta}) \leq \beta,$

with $n_{\theta} = 1 + n(n+1)/2 + nn_u + n_u + n_u(n_u + n_c)$

Extract *N* samples of the uncertain matrices according to the probability distribution **Solve**

$$\begin{split} & \gamma_{sc}^{2}(s) = \min_{\left\{\gamma^{2}, Q, Y, U, X\right\}} \gamma^{2}, \quad subject \ to \ Q = Q^{T} > 0, \ U > 0 \ diagonal, \\ & \operatorname{He} \begin{bmatrix} A_{cl}^{(i)} Q & B_{cl,u}^{(i)} U + B_{cl,v}^{(i)} X + Y^{T} & B_{cl,w}^{(i)} & 0 \\ C_{cl,u}^{(i)} Q & D_{cl,ud}^{(i)} U + D_{cl,uv}^{(i)} X - U & D_{cl,uw}^{(i)} & 0 \\ 0 & 0 & -I/2 & 0 \\ C_{cl,z}^{(i)} Q & D_{cl,zd}^{(i)} U + D_{cl,zv}^{(i)} X & D_{cl,zw}^{(i)} - \frac{\gamma^{2}}{2}I \end{bmatrix} < 0, \quad \begin{bmatrix} Q & Y_{lk}^{T} \\ Y_{lk} & \overline{u}_{k}^{2}/s^{2} \\ \forall k = 1, \dots, n_{u} \\ \forall i = 1, \dots, N \end{split}$$

If the above LMIs are feasible, select the static anti-windup gain $D_{aw} = XU^{-1}$

Then, for each $\|w\|_2 < s$, the zero initial state solution of the closed loop satisfies $Pr(\|z\|_2 > \gamma_{sc}(s) \|w\|_2) \le \varepsilon$, with probability no smaller than $1 - \beta$.

Analysis conditions can also be easily formulated

Nominal analysis and desig

Robust anti-windup

Non-common certificates

References

Static AW synthesis based on scenario with certificates

Theorem (Robust static AW using scenario with certificates Formentin et al. [2014])

Fix a positive value s $\geq \|w\|_2$, $\varepsilon \in (0, 1)$, $\beta \in (0, 1)$, and select N satisfying

 $\mathsf{B}(N,\varepsilon, \mathbf{n}_{\theta}) \leq \beta,$

with $n_{\theta} = 1 + \frac{n(n+1)}{2 + m_u} + n_u + n_u(n_u + n_c)$

Extract *N* samples of the uncertain matrices according to the probability distribution **Solve**

$$\begin{split} \gamma_{sc}^{2}(s) &= \min_{\substack{\{\gamma^{2}, U, X\}, \{Q_{i}, Y_{i}\}}} \gamma^{2}, \quad \text{subject to } Q_{i} = Q_{i}^{-T} > 0, \ U > 0 \ \text{diagonal}, \\ \mathrm{He} \begin{bmatrix} A_{cl}^{(i)}Q_{i} & B_{cl,d}^{(i)}U + B_{cl,v}^{(i)}X + Y_{i}^{T} & B_{cl,w}^{(i)} & 0\\ C_{cl,u}^{(i)}Q_{i} & D_{cl,ud}^{(i)}U + D_{cl,wv}^{(i)}X - U & D_{cl,ww}^{(i)} & 0\\ 0 & 0 & -I/2 & 0\\ C_{cl,z}^{(i)}Q_{i} & D_{cl,zd}^{(i)}U + D_{cl,zv}^{(i)}X & D_{cl,wv}^{(i)} - \frac{\gamma^{2}}{2}I \end{bmatrix} < 0, \quad \begin{bmatrix} Q_{i} & Y_{i[k]}^{T} \\ Y_{i[k]} & \overline{u}_{k}^{2}/s^{2} \\ \forall k = 1, \dots, n_{u} \\ \forall i = 1, \dots, N \end{split} \ge 0, \end{split}$$

If the above LMIs are feasible, select the static anti-windup gain $D_{aw} = XU^{-1}$

Then, for each $||w||_2 < s$, the zero initial state solution of the closed loop satisfies $Pr(||z||_2 > \gamma_{sc}(s) ||w||_2) \le \varepsilon$, with probability no smaller than $1 - \beta$.

Analysis conditions can also be easily formulated

| Illustrative ex | ample: | Another | (larger) | passive network | |
|----------------------------|------------------|--------------|--------------------|-------------------------|--------|
| 000 | 000000000 | | 0000000000 | 000000000000 | |
| Saturation and anti-windup | Nominal analysis | s and design | Robust anti-windup | Non-common certificates | Refere |



• Uncertain parameters with (known) Gaussian distribution

| parameter | mean | std dev | parameter | mean | std dev |
|----------------|-------|----------|-----------------------|--------|----------|
| R_1 | 313 Ω | \pm 10 | R_5 | 10 F | \pm 10 |
| R_2 | 20 Ω | \pm 10 | <i>C</i> ₁ | 0.01 F | \pm 10 |
| R ₃ | 315 Ω | \pm 10 | <i>C</i> ₂ | 0.01 F | \pm 10 |
| R ₄ | 17 Ω | \pm 10 | <i>C</i> 3 | 0.01 F | \pm 10 |

- Input generator voltage constrained:
 - $u(t) = V_i(t) \in [-\overline{u}, \overline{u}] = [-1 Volt, 1 Volt]$
- Design parameters are $\varepsilon = 0.01, \beta = 10^{-6}, s = 0.003, n_{\theta} = 35$ $\Rightarrow N = 2270 \text{ (not 7565) for design based on sequential algorithm}$



- Robust compensator shows better robust performance (red curves)
- The nominal behavior slightly deteriorated (thin curves)



Without anti-windup (black dashed), with nominal anti-windup (blue dashed-dotted) and with robust anti-windup (red solid)



Time responses confirm nonlinear \mathcal{L}_2 gain trends



Nominal analysis and desig

Robust anti-windup

Non-common certificates

References

Optimization of the reachable set estimate

Theorem (Robust static AW using scenario with certificates Formentin et al. [2014])

Fix a positive value s $\geq \|w\|_2, \, \varepsilon \in (0,1), \, \beta \in (0,1),$ and select N satisfying

 $\mathsf{B}(N,\varepsilon,\mathbf{n}_{\theta}) \leq \beta,$

with $n_{\theta} = n(n+1)/2 + n_u + n_u(n_u + n_c)$

Extract *N* samples of the uncertain matrices according to the probability distribution **Solve**

$$\begin{split} \gamma_{sc}^{2}(s) &= \min_{\substack{\{\bar{Q}, U, X\}, \{Q_{i}, Y_{i}\}}} \text{trace}(\bar{Q}), \quad \text{subject to } Q_{i} = Q_{i}^{T} > 0, \ U > 0 \ \text{diagonal}, \\ \bar{Q} &\geq Q_{i} \\ \text{He} \begin{bmatrix} A_{cl}^{(i)} Q_{i} & B_{cl,d}^{(i)} U + B_{cl,v}^{(i)} X + Y_{i}^{T} & B_{cl,w}^{(i)} \\ C_{cl,u}^{(i)} Q_{i} & D_{cl,ud}^{(i)} U + D_{cl,uv}^{(i)} X - U & D_{cl,uw}^{(i)} \\ 0 & 0 & -I/2 \end{bmatrix} < 0, \quad \begin{cases} Q_{i} & Y_{i[k]}^{T} \\ Y_{i[k]} & \overline{u}_{k}^{2}/s^{2} \end{bmatrix} \geq 0, \\ \forall k = 1, \dots, n_{u} \\ \forall i = 1, \dots, N \end{cases}$$
(1)

If the above LMIs are feasible, select the static anti-windup gain $D_{aw} = XU^{-1}$

Then, for each $||w||_2 < s$, the zero initial state solution of the closed loop has probability $1 - \epsilon$ of remaining in the ellipsoid $\mathcal{E}(\bar{Q}, s)$ with level of confidence no smaller than $1 - \beta$.

Analysis conditions can also be easily formulated

Nominal analysis and desig

Robust anti-windup

Non-common certificates

References

Reachable sets for simple 2D example

Left is Nominal design:

Clearly unsuitable Nominal parameters (black), Perturbed parameters (blue)

Right is Robust design:

Potential behind noncommon Q_i 's Guaranteed region $\mathcal{E}(\bar{Q}, s)$ (black), A collection of sets $\mathcal{E}(Q_i, s)$ (blue)



| Saturation and anti-windup | Nominal analysis and design | Robust anti-windup 0000000000 | Non-common certificates | References |
|----------------------------|-----------------------------|----------------------------------|-------------------------|------------|
| Concluding re | emarks | | | |

- Deterministic formulation of the robust static anti-windup design problem is nonconvex
- Scenario approach can be used for **robust static anti-windup** compensator synthesis and for robust stability and performance analysis with **common certificates**
- New tool scenario with certificates allows for non-common certificates and results with reduced conservativeness
- Current/future work:
 - transform *s* into a random variable to deal with uncertain disturbances/references
 - address robust dynamic anti-windup compensation

| Saturation and anti-windup 000 | Nominal analysis and design | Robust anti-windup 0000000000 | Non-common certificates | References |
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