Decentralized Estimation and Control Methods for Cooperative Robot Motion

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- 1. Graphs, Matrices, and Eigenvalues
- 2. Connectivity vs Infinitesimal Rigidity
- 3. Maintenance Problems and Methods
- 4. Handling Multiple Objectives in Maintenance Problems
- 5. Applications



Partial list:

- P. Yang, R.A. Freeman, G.J. Gordon, K.M. Lynch, S.S. Srinivasa, and R. Sukthankar, "Decentralized estimation and control of graph connectivity for mobile sensor networks," Automatica, vol. 46, no. 2. pp. 390–396, Feb. 2010.
- G. Hollinger and S. Singh, "Multirobot coordination with periodic connectivity: Theory and experiments," IEEE Transactions on Robotics , 2012,
- L. Sabattini, C. Secchi, N. Chopra, and A Gasparri. Distributed Control of Multirobot Systems With Global Connectivity Maintenance. Robotics, IEEE Transactions on Robotics, 29(5):1326-1332, 2013.
- D. Carboni, R.K. Williams, A. Gasparri, G. Ulivi, and G.S. Sukhatme. Rigidity-Preserving Team Partitions in Multi-Agent Networks. IEEE Transactions on Cybernetics, pp 1-14, 2014.



If you want to know more about what follows:

- Robuffo Giordano, P., A. Franchi, C. Secchi, and H. H. Bülthoff (2013). *A Passivity-Based Decentralized Strategy for Generalized Connectivity* Maintenance". The International Journal of Robotics Research 32.3, pp. 299–323.
- Zelazo, D., A. Franchi, H. H. Bülthoff, and P. Robuffo Giordano (2014). Decentralized Rigidity Maintenance Control with Range Measurements for Multi-Robot Systems. The International Journal of Robotics Research 34.1, pp. 105–128.
- Nestmeyer T., P. Robuffo Giordano, H. H. Bülthoff, and A. Franchi (2016), Decentralized Simultaneous Multi-target Exploration using a Connected Network of Multiple Robots. Autonomous Robots, online first June 2016, doi:10.1007/s10514-016-9578-9

Graphs, Matrices, and Eigenvalues

Graph



- $\mathcal{G} = (\mathcal{V}, \, \mathcal{E})$ is an undirected graph or simply graph
 - $\mathcal{V} = \{1, \dots, N\}$ vertex set
 - $\mathcal{E} \subset (\mathcal{V} \times \mathcal{V}) / \sim$ edge set
 - \sim equivalence relation identifying (i, j) and (j, i)



A Graph models an Adjacency Structure

 $[(i, j)] \in \mathcal{E} \Leftrightarrow$ vertexes *i* and *j* are **neighbors** or **adjacent**

• (i, j), i < j representative element of the equivalence class [(i, j)]

$$[\mathcal{V} \times \mathcal{V}] = \{(1, 2), (1, 3), \dots, (1, N), \dots, (N-1, N)\} \\ = \{e_1, e_2, \dots, e_{N-1}, \dots, e_{N(N-1)/2}\}$$

- $[(i, i)] \notin \mathcal{E}, \forall i \in \mathcal{V} \text{ (no self-loops)}$
- $\mathcal{N}_i = \{j \in \mathcal{V} \mid (i, j) \in \mathcal{E}\}$ set of **neighbors** of *i*



 $E \in \mathbb{R}^{N \times N(N-1)/2}$ is the (full) incidence matrix of \mathcal{G}

- $\forall e_k = (i, j) \in [\mathcal{V} \times \mathcal{V}]:$
 - $E_{ik} = -1$ and $E_{jk} = 1$, if $e_k \in \mathcal{E}$
 - $E_{ik} = 0$ and $E_{jk} = 0$, otherwise

Matricial representation of a graph

Example:



$$E = \begin{pmatrix} 1 & 1 & 1 & 0 & 0 & 0 \\ -1 & 0 & 0 & 1 & 0 & 0 \\ 0 & -1 & 0 & -1 & 0 & 1 \\ 0 & 0 & -1 & 0 & 0 & -1 \end{pmatrix}$$
$$e_1 \quad e_2 \quad e_3 \quad e_4 \quad e_5 \quad e_6$$

remember:

$$\{e_1, e_2, \ldots, e_{N-1}, \ldots, e_{N(N-1)/2}\} = \{(1, 2), (1, 3), \ldots, (1, N), \ldots (N-1, N)\}$$

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Assume N mobile robots moving in an environment:

- $\mathbf{x}_i \in \mathbb{R}^{N_{\mathbf{x}}}$ *i*-th robot configuration, $i \in 1 \dots N$
- $\boldsymbol{z} \in \mathbb{R}^{N_{\boldsymbol{z}}}$ environment configuration

Consider two maps

 $\begin{array}{rll} \mbox{robot map} & \mbox{v}: & \ensuremath{\mathbb{R}}^{N_x} \ensuremath{\; \ni \;} x_i \ensuremath{\; \mapsto \;} v(x_i) = \mbox{v}_i \in \ensuremath{\mathbb{R}}^{N_v} \\ \mbox{connection map} & \mbox{w}: & \ensuremath{\mathbb{R}}^{N_x} \times \ensuremath{\mathbb{R}}^{N_z} \ensuremath{\; \ni \;} (x_i, x_j, z) \ensuremath{\; \mapsto \;} w(x_i, x_j, z) = \mbox{w}_{ij} \in \ensuremath{\mathbb{R}}_{\geq 0} \end{array}$

with the properties

•
$$\mathbf{w}_{ij} = \mathbf{w}_{ji}$$
 (symmetry)

• **w**_{ii} = 0

example: what can those maps model?



The connection map w defines an associated graph $\mathcal{G} = (\mathcal{V}, \mathcal{E})$, where

- $\mathcal{V} = \{1, 2, \dots, N\}$
- $\mathcal{E} = \{ e_k = (i, j) \mid \mathbf{w}_{ij} > 0 \}$
- the **positive weight** \mathbf{w}_{ij} is associated to each edge $(i, j) \in \mathcal{E}$

Both maps **v** and **w** define an associated framework $(\mathcal{G}, \mathbf{v})$ where

- ${\mathcal G}$ is the associated graph
- \mathbf{v}_i is associated to each vertex $i \in \mathcal{V}$



$$A = \begin{pmatrix} 0 & \mathbf{w}_{12} & \dots & \mathbf{w}_{1N} \\ \mathbf{w}_{12} & 0 & \dots & \mathbf{w}_{2N} \\ \vdots & \vdots & \ddots & \vdots \\ \mathbf{w}_{1N} & \mathbf{w}_{2N} & \cdots & 0 \end{pmatrix}$$

 $\in \mathbb{R}^{\textit{N} \times \textit{N}}$ is the adjacency (or weight) matrix of $\mathcal G$

Note that

- $A_{ij} = 0$ if $(i, j) \notin \mathcal{E}$
- $A_{ij} > 0$ otherwise

Properties:

$$\mathsf{P.1} \ \mathsf{A} = \mathsf{A}(\mathsf{x}_1, \ldots, \mathsf{x}_N, \mathsf{z})$$

P.2 A is square

$$\begin{array}{l} \mathsf{P.3} \ A_{ij} = A_{ij} \ (\text{symmetric}) \\ \mathsf{P.4} \ A_{ij} = A_{ij} \geq 0 \ (\text{nonnegative}) \\ \mathsf{P.5} \ A_{ii} = 0 \end{array}$$

Example:



$$A = \begin{pmatrix} 0 & \mathbf{w}_{12} & \mathbf{w}_{13} & \mathbf{w}_{14} \\ \mathbf{w}_{12} & 0 & \mathbf{w}_{23} & 0 \\ \mathbf{w}_{13} & \mathbf{w}_{23} & 0 & \mathbf{w}_{34} \\ \mathbf{w}_{14} & 0 & \mathbf{w}_{34} & 0 \end{pmatrix}$$



$$L = \begin{pmatrix} \sum_{j=1}^{n} \mathbf{w}_{1j} & -\mathbf{w}_{12} & \dots & -\mathbf{w}_{1N} \\ -\mathbf{w}_{12} & \sum_{j=1}^{n} \mathbf{w}_{j2} & \dots & -\mathbf{w}_{2N} \\ \vdots & \vdots & \ddots & \vdots \\ -\mathbf{w}_{1N} & -\mathbf{w}_{2N} & \dots & \sum_{j=1}^{n} \mathbf{w}_{jN} \end{pmatrix}$$

 $\in \mathbb{R}^{\textit{N} \times \textit{N}}$ is the Laplacian matrix of $\mathcal G$

Note that

Example:

•
$$L = \operatorname{diag}(\delta_i) - A$$
,

where
$$\delta_i = \sum_{j=1}^n \mathbf{w}_{ij}$$

(degree of vertex *i*)

1 (1,3) 3 (3,q) (1,4) 4

Properties:

P.1
$$L = L(\mathbf{x}_1, \dots, \mathbf{x}_N, \mathbf{z})$$

P.2 L is square
P.3 $L_{ii} = L_{ii}$ (symmetric)

$$L = \begin{pmatrix} \mathbf{w}_{12} + \mathbf{w}_{13} + \mathbf{w}_{14} & -\mathbf{w}_{12} & -\mathbf{w}_{13} & -\mathbf{w}_{14} \\ -\mathbf{w}_{12} & \mathbf{w}_{12} + \mathbf{w}_{23} & -\mathbf{w}_{23} & \mathbf{0} \\ -\mathbf{w}_{13} & -\mathbf{w}_{23} & \mathbf{w}_{13} + \mathbf{w}_{23} + \mathbf{w}_{34} & -\mathbf{w}_{34} \\ -\mathbf{w}_{14} & \mathbf{0} & -\mathbf{w}_{34} & \mathbf{w}_{14} + \mathbf{w}_{34} \end{pmatrix}$$



Connectivity

 ${\mathcal G}$ is connected if there is a path between every pair of vertices, i.e.,

 $\forall i \in \mathcal{V} \text{ and } j \in \mathcal{V} ackslash i, \quad \exists \text{ a path (sequence of adjacent edges) from } i \text{ to } j$

This is a combinatorial definition of connectivity



disconnected graph

connected graph

question: connectivity is a global property, what does it mean? and why it is global?



What connectivity can model?

- connected communication network
- connected **sensing** network
- connected **control** network
- connected **planning** roadmap

What connectivity is important for?

- pass a message from any robot to any other robot
- know the relative position between any two robots in a common frame
- converge to a common point
- share a common goal

Related concepts

- group, cohesiveness
- aggregation
- sharing



Additional properties of $L = \operatorname{diag}(\delta_i) - A$

• L is positive semi-definite, i.e., all the eigenvalues are real and non-negative

$$0 \leq \lambda_1 \leq \lambda_2 \leq \ldots \leq \lambda_N$$

• $\sum_{i=1}^{n} L_{ij} = 0$ $\forall i = 1 \dots N$, i.e., $L\mathbf{1} = \mathbf{0}$, therefore

 $\lambda_1 = 0$ and it is associated to the eigenvector $\mathbf{1} = \begin{pmatrix} 1 & 1 & \dots & 1 \end{pmatrix}^T$

(Fiedler 1973)

 $\lambda_2 > 0$ if the graph ${\cal G}$ is connected and $\lambda_2 = 0$ otherwise



 λ_2 provides an algebraic definition of connectivity

 $\Rightarrow \lambda_2$ is called *algebraic connectivity, connectivity eigenvalue*, or Fiedler eigenvalue $\lambda_2 = \lambda_2(\mathbf{x}_1, \dots, \mathbf{x}_N, \mathbf{z})$ is a global quantity

Example (if $\mathbf{w}_{ij} \in \{0,1\} \ \forall (i,j) \in \mathcal{V} \times \mathcal{V}$):





A framework of positions is a particular framework $(\mathcal{G}, \mathbf{v})$ in the special case in which $\mathbf{v} : \mathcal{V} \to \mathbb{R}^d$ maps each vertex to the position in \mathbb{R}^d of the *i*-th robot

• if
$$d = 2$$
, $\mathbf{v}_i = \mathbf{p}_i = \begin{pmatrix} p_i^x \\ p_i^y \end{pmatrix}$, 2D position of robot i
• if $d = 3$, $\mathbf{v}_i = \mathbf{p}_i = \begin{pmatrix} p_i^x \\ p_i^y \\ p_i^z \end{pmatrix}$, 3D position of robot i

In the following

- it will be (mainly) d = 3, similar results apply for d = 2
- · we refer only to framework of positions, called simply frameworks

Equivalent and Congruent Frameworks



Consider two frameworks $(\mathcal{G}, \boldsymbol{p}')$ and $(\mathcal{G}, \boldsymbol{p}'')$

- same graph ${\mathcal G}$
- different positions \boldsymbol{p}' and \boldsymbol{p}''

Frameworks $(\mathcal{G}, \mathbf{p}')$ and $(\mathcal{G}, \mathbf{p}'')$ are

- equivalent: if $\|\mathbf{p}'_i \mathbf{p}'_j\| = \|\mathbf{p}''_i \mathbf{p}''_j\|$ for all $(i, j) \in \mathcal{E}$, and
- congruent: if $\|\mathbf{p}'_i \mathbf{p}'_j\| = \|\mathbf{p}''_i \mathbf{p}''_j\|$ for all $(i, j) \in \mathcal{V} \times \mathcal{V}$





Global Rigidity

The framework (\mathcal{G},p') is globally rigid if every other framework (\mathcal{G},p'') which

- is equivalent to $(\mathcal{G}, \mathbf{p}'')$
- is also congruent to $(\mathcal{G},\mathbf{p}')$

This is, again, a combinatorial definition



Rigidity



Rigidity

The framework $(\mathcal{G}, \mathbf{p}')$ is **rigid** if $\exists \epsilon > 0$ such that every other framework $(\mathcal{G}, \mathbf{p}'')$ which

- is equivalent to $(\mathcal{G}, \boldsymbol{p}'')$ and
- satisfies $\|\mathbf{p}'_i \mathbf{p}''_i\| < \epsilon$ for all $i \in \mathcal{V}$,

is congruent to $(\mathcal{G}, \boldsymbol{p}')$

This is, again, a combinatorial definition



question: is rigidity a global property of the graph as well?



What rigidity can model?

• rigid mechanical structure made of bars

but also:

- rigid sensing network
- rigid control network

What rigidity is important for?

- **univocally** compute the arrangement (**shape**) of a group of robots only using **inter-distances**
- achieve (or track) a desired shape **only controlling** the **inter-distances** (formation control)

Related concepts

- parallel rigidity
- persistent graph
- tensegrity



question: do you know an example of use of rigidity in robotics?



question: do you know an example of use of rigidity in robotics?

6-DOF Stewart platform parallel robot



Credits: Robert L. Williams II



Let's give a definition of rigidity that is differential (\Leftrightarrow involves infinitesimal motions)

Consider a trajectory $\mathbf{p}(t)$ with $t \ge t_0$ and impose equivalence along the trajectory:

$$\left\|\mathbf{p}_i(t)-\mathbf{p}_j(t)
ight\|^2=\left\|\mathbf{p}_i(t_0)-\mathbf{p}_j(t_0)
ight\|^2= ext{const} \quad ext{for all} \quad (i,j)\in\mathcal{E}, \ orall t\geq t_0$$

Differentiating with respect to time the constraint above:

$$\left(\mathbf{p}_{i}(t)-\mathbf{p}_{j}(t)\right)^{T}\left(\dot{\mathbf{p}}_{i}(t)-\dot{\mathbf{p}}_{j}(t)
ight)=0 \quad \text{for all} \quad (i,j)\in\mathcal{E}, \ \forall t\geq t_{0} \tag{1}$$

Trivial Motion

A collective motion that consists of only **global roto-translations** of the whole set of positions in the framework

Infinitesimal Rigidity

The framework $(\mathcal{G}, \mathbf{p}(t_0))$ is **infinitesimally rigid** if every possible motion that satisfies (1) is **trivial**



question: is this framework rigid in $\mathbb{R}^2 ?$ is it infinitesimally rigid?





question: is this framework rigid in \mathbb{R}^2 ? is it infinitesimally rigid?



- infinitesimal rigidity \Rightarrow rigidity
- rigidity \Rightarrow infinitesimal rigidity



Let us write the infinitesimal rigidity constraint in a matricial form

$$0 = \mathbf{w}_{ij} \left(\mathbf{p}_{i}(t) - \mathbf{p}_{j}(t) \right)^{T} \left(\dot{\mathbf{p}}_{i}(t) - \dot{\mathbf{p}}_{j}(t) \right) =$$

$$= \mathbf{w}_{ij} \left(\mathbf{p}_{i}(t) - \mathbf{p}_{j}(t) \right)^{T} \dot{\mathbf{p}}_{i}(t) - \left(\mathbf{p}_{i}(t) - \mathbf{p}_{j}(t) \right)^{T} \dot{\mathbf{p}}_{j}(t) =$$

$$= \mathbf{w}_{ij} \underbrace{ \left(-\mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{j}(t) \right)^{T}}_{\text{vertex } i} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{j}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\text{vertex } j} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} - \mathbf{0}^{T} - \underbrace{\left(\mathbf{p}_{i}(t) - \mathbf{p}_{i}(t) \right)^{T}}_{\mathbf{k}_{ij} \in \mathbb{R}^{1 \times 3N}} -$$

where
$$\mathbf{0} = \begin{pmatrix} 0 & 0 & \dots & 0 \end{pmatrix}^T$$



stacking the previous constraints for every $(i, j) \in \{e_1, e_2 \dots e_{N-1} \dots, e_{N(N-1)/2}\}$:

$$\underbrace{\begin{pmatrix} \mathbf{w}_{12} \\ \vdots \\ \mathbf{w}_{N(N-1)} \end{pmatrix}}_{W(\mathbf{w}) \in \mathbb{R}^{\frac{N(N-1)}{2} \times \frac{N(N-1)}{2}} \times \frac{N(N-1)}{2}} \underbrace{\begin{pmatrix} \mathbf{k}_{12} \\ \vdots \\ \mathbf{k}_{N(N-1)} \end{pmatrix}}_{K(\mathbf{p})} \underbrace{\begin{pmatrix} \dot{\mathbf{p}}_{1} \\ \vdots \\ \dot{\mathbf{p}}_{N} \end{pmatrix}}_{\dot{\mathbf{p}} \in \mathbb{R}^{3N}} = \underbrace{W(\mathbf{w}) K(\mathbf{p})}_{R(\mathbf{w},\mathbf{p}) \in \mathbb{R}^{\frac{N(N-1)}{2} \times 3N}} \dot{\mathbf{p}} = R(\mathbf{w},\mathbf{p}) \dot{\mathbf{p}} = \mathbf{0}$$

Rigidity Matrix

 $R(\mathbf{w}, \mathbf{p})$ is the (weighted) rigidity matrix





$$d = 2 (\mathbb{R}^2)$$

 $N = 4$
 $N(N-1)/2 = 6$

$$\begin{split} R(\mathbf{w},\mathbf{p}) = \\ & \left(\begin{array}{c} \mathsf{w}_{12}(p_1^x - p_2^x) \; \mathsf{w}_{12}(p_1^y - p_2^y) \; \; \mathsf{w}_{12}(p_2^x - p_1^x) \; \mathsf{w}_{12}(p_2^y - p_1^y) \; & 0 \; & 0 \; & 0 \; \\ \mathsf{w}_{13}(p_1^x - p_3^x) \; \mathsf{w}_{13}(p_1^y - p_3^y) \; & 0 \; & 0 \; & \mathsf{w}_{13}(p_3^x - p_1^x) \; \mathsf{w}_{13}(p_3^y - p_1^y) \; & 0 \; & 0 \; \\ \mathsf{w}_{14}(p_1^x - p_4^x) \; \mathsf{w}_{14}(p_1^y - p_4^y) \; & 0 \; & 0 \; & 0 \; & \mathsf{w}_{14}(p_3^x - p_1^x) \; \mathsf{w}_{14}(p_4^y - p_1^y) \; \\ \mathsf{0} \; & \mathsf{0} \; \; \mathsf{w}_{23}(p_2^x - p_3^x) \; \mathsf{w}_{23}(p_2^y - p_3^y) \; \mathsf{w}_{23}(p_3^x - p_2^x) \; \mathsf{w}_{23}(p_3^y - p_2^y) \; & \mathsf{0} \; & \mathsf{0} \; \\ \mathsf{0} \; \; \mathsf{0} \; \; \mathsf{0} \; \; \mathsf{0} \; \; \mathsf{0} \; \mathsf{0} \; & \mathsf{0} \; & \mathsf{0} \; \\ \mathsf{0} \; \; \mathsf{0} \; \; \mathsf{0} \; \; \mathsf{0} \; \\ \mathsf{0} \; \; \mathsf{0} \;$$



- rigidity is defined combinatorially ("...s.t. every other framework...")
- infinitesimal rigidity implies rigidity
- converse not true (degenerate cases) but...
- infinitesimal rigidity can be defined algebraically, in fact...



- collective roto-translations in \mathbb{R}^3 keep constant all the distances, by definition, i.e., if $\dot{\mathbf{p}}$ is trivial then $R(\mathbf{w}, \mathbf{p})\dot{\mathbf{p}} = 0$
- \Rightarrow Dim (ker[$R(\mathbf{w}, \mathbf{p})$]) \geq 6 always
- for infinitesimally rigid frameworks the motion that keep constant all the distances are only collective roto-translations in ℝ³
 i.e., if R(w, p)p = 0 then p is trivial
- infinitesimally rigidity \Rightarrow Dim (ker[$R(\mathbf{w}, \mathbf{p})$]) = 6

(Tay and Whiteley 1985) and (Zelazo et al. 2014)

A framework is infinitesimally rigid if and only if $rank[R(\mathbf{w}, \mathbf{p})] = 3N - 6$

• despite its name, the rigidity matrix is actually characterizing **infinitesimal rigidity** (rather than **rigidity**)



 $S(\mathbf{w}, \mathbf{p}) = R(\mathbf{w}, \mathbf{p})^T R(\mathbf{w}, \mathbf{p}) \in \mathbb{R}^{3N \times 3N}$ is the symmetric rigidity matrix

(Zelazo et al. 2014)

Properties:

- $\mathsf{P.1} \ S = S(\mathbf{w}, \mathbf{p}) = S(\mathbf{x}_1, \dots, \mathbf{x}_N, \mathbf{z})$
- P.2 $S \in \mathbb{R}^{3N \times 3N}$ (square)
- P.3 $S_{ij} = S_{ji}$ (symmetric)
- P.4 Dim $(\ker[S(\mathbf{w}, \mathbf{p})]) \ge 6$

(Zelazo et al. 2014)

A framework is infinitesimally rigid if and only if $rank[S(\mathbf{w}, \mathbf{p})] = 3N - 6$

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Additional properties of $S = R^T R$

• S is positive semi-definite, i.e., all the eigenvalues are real and non-negative

 $0 \leq \varsigma_1 \leq \varsigma_2 \leq \ldots \leq \varsigma_6 \leq \varsigma_7 \leq \ldots \leq \varsigma_{3N}$

• $Dim(ker[S(\mathbf{w}, \mathbf{p})]) \ge 6$, therefore

$$\varsigma_1 = \varsigma_2 = \varsigma_3 = \varsigma_4 = \varsigma_5 = \varsigma_6 = 0$$

(Zelazo et al. 2014)

 $\varsigma_7>0$ if the framework is infinitesimally rigid and $\varsigma_7=0$ otherwise

 ς_7 provides an algebraic definition of infinitesimal rigidity $\Rightarrow \varsigma_7$ is called the **rigidity eigenvalue** (Zelazo et al. 2014) $\varsigma_7 = \varsigma_7(\mathbf{x}_1, \dots, \mathbf{x}_N, \mathbf{z})$ is a global quantity

Connectivity vs Infinitesimal Rigidity

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Connectivity

 \exists a path between any pair of vertexes

- depends on x₁,..., x_N, z
 (global property)
- Laplacian matrix $L \in \mathbb{R}^{N \times N}$
- \Leftrightarrow Fidler eigenvalue $\lambda_2 > 0$

Infinitesimal rigidity

distance-preservation on the edges forces a trivial (roto-translational) movement

- depends on x₁,..., x_N, z
 (global property)
- symmetric rigidity matrix $S \in \mathbb{R}^{3N imes 3N}$
- \Leftrightarrow rigidity eigenvalue $\varsigma_7 > 0$



(Infinitesimal) Rigidity \Rightarrow Connectivity, i.e., $\varsigma_7 > 0 \Rightarrow \lambda_2 > 0$

In fact, e.g., by contradiction:

- not connected implies at least two connected components
- **distance** between the two connected components **can change** still preserving equivalence
- \Rightarrow by enforcing infinitesimal rigidity one enforces connectivity as well


Connectivity

- applicable to any graph
- depends only on w
- \Rightarrow infinitesimal rigidity

Infinitesimal rigidity

- applicable only to frameworks (graphs + positions)
- depends both on w and v=p
- \Rightarrow connectivity

Infinitesimal rigidity is a stronger property and applies to a more particular structure (framework)

Maintenance Problems and Methods



Assume each robot $i = 1, \ldots, N$

- can control $\mathbf{x}_i(t)$, $\forall t \geq t_0$ (with $\mathbf{x}_i(t)$ smooth enough)
- has some objectives (mission)

Maintenance problem(s)

Maintenance \neq

- assume G is connected (or (G, \mathbf{p}) is infinitesimally rigid) for $t = t_0$
- control $\mathbf{x}_1(t), \ldots, \mathbf{x}_N(t)$ such that
 - 1. \mathcal{G} stays connected (or $(\mathcal{G}, \mathbf{p})$ stays infinitesimally rigid) $\forall t > t_0$
 - 2. the mission of each robot is accomplished
 - eventual achievement
 - periodical achievement



Using the algebraic formulation of connectivity and infinitesimal rigidity

Connectivity maintenance

- assume λ₂(t₀) > 0
- for *t* > *t*₀
 - maintain $\lambda_2(\mathbf{x}_1(t), \ldots, \mathbf{x}_N(t), \mathbf{z}) > 0$
 - o and accomplish the mission

Infinitesimal rigidity maintenance

- assume
 _{γ7}(t₀) > 0
- for $t > t_0$
 - maintain $\varsigma_7(\mathbf{x}_1(t), \dots, \mathbf{x}_N(t), \mathbf{z}) > 0$
 - and accomplish the mission





- 1. define potential function $V:(\mu_{\min},+\infty)
 ightarrow\mathbb{R}^+$, that
 - $\begin{array}{l} \circ \ \mbox{grows unbounded as } \mu \rightarrow^+ \mu^{\min} > 0 \\ \circ \ \mbox{vanishes (with vanishing derivatives) as} \\ \mu > \mu^0 > \mu^{\min} \\ \circ \ \mbox{is, at least, } C^1, \ \mbox{i.e., it exists } \frac{dV}{d\mu}, \ \forall \mu > \mu^{\min} \end{array}$



2. let each robot command

$$\mathbf{x}_{i}^{(h)} = \left. \frac{\mathrm{d}V}{\mathrm{d}\mu} \right|_{\lambda_{2}(t)} \left. \frac{\partial\lambda_{2}}{\partial\mathbf{x}_{i}} \right|_{(\mathbf{x}_{1},\ldots,\mathbf{x}_{N},\mathbf{z})} + u_{i}$$

(for connectivity maintenance)

$$\mathbf{x}_{i}^{(h)} = \left. \frac{\mathrm{d}V}{\mathrm{d}\mu} \right|_{\varsigma_{7}(t)} \left. \frac{\partial\varsigma_{7}}{\partial \mathbf{x}_{i}} \right|_{(\mathbf{x}_{1},\ldots,\mathbf{x}_{N},\mathbf{z})} + u_{i}$$

(for infinitesimal rigidity maintenance)

where u_i is a properly designed additional control input accounting for

- accomplishment of mission
- stability



connectivity maintenance

$$\frac{\mathrm{d}V}{\mathrm{d}\mu}\bigg|_{\lambda_2(t)} \left.\frac{\partial\lambda_2}{\partial\mathbf{x}_i}\right|_{(\mathbf{x}_1,\ldots,\mathbf{x}_N,\mathbf{z}_N)}$$

infinitesimal rigidity maintenance

$$\frac{\mathrm{d}\boldsymbol{V}}{\mathrm{d}\boldsymbol{\mu}}\bigg|_{\varsigma_{7}(t)}\left.\frac{\partial\varsigma_{7}}{\partial\mathbf{x}_{i}}\right|_{(\mathbf{x}_{1},\ldots,\mathbf{x}_{N},\mathbf{z})}$$

Gradient computation is composed by two parts



First part: computation of
$$\frac{\mathrm{d}V}{\mathrm{d}\mu}\Big|_{\lambda_2(t)}$$
 (or $\frac{\mathrm{d}V}{\mathrm{d}\mu}\Big|_{\varsigma_7(t)}$)

requires that each robot knows:

- \bullet the function V
- λ₂(t) (or ς₇(t))



Second part: Computation of $\left. \frac{\partial \lambda_2}{\partial \mathbf{x}_i} \right|_{(\mathbf{x}_1,...,\mathbf{x}_N,\mathbf{z})}$ (or $\left. \frac{\partial \varsigma_7}{\partial \mathbf{x}_i} \right|_{(\mathbf{x}_1,...,\mathbf{x}_N,\mathbf{z})}$)

requires in general

• the analytic expression of the gradient of λ_2 (or ς_7) with respect to \mathbf{x}_i

Gradient of λ_2 and ς_7



Given a matrix M, any eigenvalue can be written as $\mu = \mathbf{u}^T M \mathbf{u}$, where

• **u** is a normalized eigenvector associated to μ (i.e., $M\mathbf{u} = \mu\mathbf{u}$ and $\mathbf{u}^T\mathbf{u} = 1$)

Connectivity

$$\lambda_2 = \mathbf{u}^T L \mathbf{u}$$

differentiating, we obtain (Yang et al. 2010)

$$\frac{\partial \lambda_2}{\partial \mathbf{x}_i} = \sum_{(j,h) \in \mathcal{E}} \frac{\partial \mathbf{w}_{jh}}{\partial \mathbf{x}_i} (\mathbf{u}_j - \mathbf{u}_h)^2$$

Infinitesimal rigidity

$$\varsigma_7 = \mathbf{u}^T S \mathbf{u}$$

differentiating, we obtain (Zelazo et al. 2014)

$$\begin{aligned} \frac{\partial \varsigma_7}{\partial \mathbf{x}_i} &= \sum_{(j,h)\in\mathcal{E}} \frac{\partial \mathbf{w}_{jh}}{\partial \mathbf{x}_i} \mathbf{s}_{jh} + \frac{\partial \mathbf{s}_{jh}}{\partial \mathbf{x}_i} \mathbf{w}_{jh} \\ \mathbf{s}_{jh} &= \left((p_j^x - p_h^x)^2 (\mathbf{u}_j^x - \mathbf{u}_h^x)^2 + (p_j^y - p_h^y)^2 (\mathbf{u}_j^y - \mathbf{u}_h^y)^2 + (p_j^z - p_h^z)^2 (\mathbf{u}_j^z - \mathbf{u}_h^z)^2 + 2(p_j^x - p_h^z) (p_j^y - p_h^y) (\mathbf{u}_j^x - \mathbf{u}_h^z) (\mathbf{u}_j^y - \mathbf{u}_h^y) + 2(p_j^x - p_h^z) (p_j^z - p_h^z) (\mathbf{u}_j^x - \mathbf{u}_h^z) (\mathbf{u}_j^z - \mathbf{u}_h^z) + 2(p_j^y - p_h^y) (p_j^z - p_h^z) (\mathbf{u}_j^x - \mathbf{u}_h^z) (\mathbf{u}_j^z - \mathbf{u}_h^z) + 2(p_j^y - p_h^y) (p_j^z - p_h^z) (\mathbf{u}_j^y - \mathbf{u}_h^y) (\mathbf{u}_j^z - \mathbf{u}_h^z) (\mathbf{u}_j^z - \mathbf{u}_h^z) \end{aligned}$$



Decentralized control law

Consider a network of robots performing a control law

The control law is decentralized if, for each robot *i*, the size of the

- communication bandwidth
- computation time (per step)
- memory used (inputs, outputs, local variables)

depends only on $|\mathcal{N}_i|$ and not on N

• a control law that is not decentralized is not scalable

Example of decentralized control law: consensus

$$\dot{\mathbf{x}}_i = \sum_{j \in \mathcal{N}_i} (\mathbf{x}_j - \mathbf{x}_i) \quad \forall i$$



The two control laws shown so far, i.e.,

connectivity maintenance

$$\frac{\mathrm{d}V}{\mathrm{d}\mu}\bigg|_{\mu=\lambda_2}\sum_{(j,h)\in\mathcal{E}}\frac{\partial\mathbf{w}_{jh}}{\partial\mathbf{x}_i}(\mathbf{u}_j-\mathbf{u}_h)^2$$

infinitesimal rigidity maintenance

$$\frac{\mathrm{d}V}{\mathrm{d}\mu}\bigg|_{\mu=\varsigma_{7}}\sum_{(j,h)\in\mathcal{E}}\frac{\partial\mathbf{w}_{jh}}{\partial\mathbf{x}_{i}}\mathbf{s}_{jh}+\frac{\partial\mathbf{s}_{jh}}{\partial\mathbf{x}_{i}}\mathbf{w}_{jh}$$

are they decentralized control law?



The two control laws shown so far, i.e.,

connectivity maintenance

$$\frac{\mathrm{d}V}{\mathrm{d}\mu}\bigg|_{\mu=\lambda_2}\sum_{(j,h)\in\mathcal{E}}\frac{\partial \mathbf{w}_{jh}}{\partial \mathbf{x}_i}(\mathbf{u}_j-\mathbf{u}_h)^2\bigg|$$

infinitesimal rigidity maintenance

$$\frac{\mathrm{d}V}{\mathrm{d}\mu}\bigg|_{\mu=\varsigma_{7}}\sum_{(j,h)\in\mathcal{E}}\frac{\partial\mathbf{w}_{jh}}{\partial\mathbf{x}_{i}}\mathbf{s}_{jh}+\frac{\partial\mathbf{s}_{jh}}{\partial\mathbf{x}_{i}}\mathbf{w}_{jh}$$

They are not decentralized control law because

- each robot must know λ_2 (or ς_7) that depends on $\mathbf{x}_1(t), \ldots, \mathbf{x}_N(t), \mathbf{z}$
- each robot must know \mathbf{w}_{jh} and \mathbf{s}_{jh} , $\forall (j, h) \in \mathcal{E}$, and $\mathbf{u}_1, \ldots, \mathbf{u}_N$ that also depend on $\mathbf{x}_1(t), \ldots, \mathbf{x}_N(t), \mathbf{z}$

Goal: make the control law decentralized

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Locality assumption for the connection map w

$$\forall i \in \mathcal{V}, \forall (j, h) \in \mathcal{E} \quad \frac{\partial \mathbf{w}_{jh}}{\partial \mathbf{x}_i} = 0 \text{ if neither } j = i \text{ nor } h = i$$

Consequence for connectivity gradient

$$\begin{split} \frac{\partial \lambda_2}{\partial \mathbf{x}_i} &= \sum_{(j,h) \in \mathcal{E}} \frac{\partial \mathbf{w}_{jh}}{\partial \mathbf{x}_i} (\mathbf{u}_j - \mathbf{u}_h)^2 = \sum_{j \in \mathcal{N}_i} \frac{\partial \mathbf{w}_{ij}}{\partial \mathbf{x}_i} (\mathbf{u}_i - \mathbf{u}_j)^2 \\ \frac{\partial \lambda_2}{\partial \mathbf{x}_i} &= \sum_{j \in \mathcal{N}_i} \mathbf{f}_\lambda \left(\frac{\partial \mathbf{w}_{ij}}{\partial \mathbf{x}_i}, \mathbf{w}_{ij}, \mathbf{x}_i, \mathbf{x}_j, \mathbf{u}_i, \mathbf{u}_j \right) \end{split}$$

Locality assumption for the connection map w

$$\forall i \in \mathcal{V}, \forall (j, h) \in \mathcal{E} \quad \frac{\partial \mathbf{w}_{jh}}{\partial \mathbf{x}_i} = 0 \text{ if neither } j = i \text{ nor } h = i$$

Consequence for infinitesimal rigidity gradient

 $\overline{j \in \mathcal{N}_i}$

where $p_{ii} = p_i - p_i$

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Locality assumption for the connection map w

 $\forall i \in \mathcal{V}, \forall (j, h) \in \mathcal{E} \quad \frac{\partial \mathbf{w}_{jh}}{\partial \mathbf{x}_i} = 0 \text{ if neither } j = i \text{ nor } h = i$

The two gradient-based control laws with locality assumption

connectivity maintenance

infinitesimal rigidity maintenance

$$V'(\lambda_2)\sum_{j\in\mathcal{N}_i}\mathbf{f}_{\lambda}\left(\frac{\partial\mathbf{w}_{ij}}{\partial\mathbf{x}_i},\mathbf{w}_{ij},\mathbf{x}_i,\mathbf{x}_j,\mathbf{u}_i,\mathbf{u}_j\right) \qquad V'(\varsigma_7)\sum_{j\in\mathcal{N}_i}\mathbf{f}_{\varsigma}\left(\frac{\partial\mathbf{w}_{ij}}{\partial\mathbf{x}_i},\mathbf{w}_{ij},\mathbf{x}_i,\mathbf{x}_j,\mathbf{u}_i,\mathbf{u}_j\right)$$

become **partially decentralized** control law, each robot must know:

• λ_2 (or ς_7) that depends on $\mathbf{x}_1(t), \ldots, \mathbf{x}_N(t), \mathbf{z}$ (not decentralized)

•
$$\mathbf{x}_i$$
, \mathbf{w}_{ij} , $\frac{\partial \mathbf{w}_{ij}}{\partial \mathbf{x}_i}$, and \mathbf{x}_j , $\forall j \in \mathcal{N}_i$, and \mathbf{z} , (decentralized)

• \mathbf{u}_i and \mathbf{u}_j , $\forall j \in \mathcal{N}_i$ that depend on $\mathbf{x}_1(t), \dots, \mathbf{x}_N(t), \mathbf{z}$ (not decentralized)

Goal: compute λ_2 (or ς_7), \mathbf{u}_i and $\mathbf{u}_j \ \forall j \in \mathcal{N}_i$ in a decentralized way



Continuous power iteration method (Yang et al. 2010; Zelazo et al. 2014)

An iterative algorithm to get an estimate $\hat{\mu}$ and $\hat{\mathbf{u}}$ of the the *I*-th eigenvalue μ and the associated eigenvector \mathbf{u} of a positive semidefinite matrix $M \in \mathbb{R}^n$

Denote with $T \in \mathbb{R}^{n \times l - 1}$ the image matrix of the first l - 1 eigenvectors

$$\dot{\hat{\mathbf{u}}} = -k_1 T T^T \hat{\mathbf{u}} - k_2 M \hat{\mathbf{u}} - k_3 \left(\frac{\hat{\mathbf{u}}^T \hat{\mathbf{u}}}{n} - 1 \right)$$

- $-k_1 T T^T \hat{\mathbf{u}}$: deflation: to remove the components spanned by the first l-1 eigenvectors
- $-k_2 M \hat{\mathbf{u}}$: direction update, to move towards \mathbf{u}
- $-k_3\left(\frac{\hat{\mathbf{u}}^T\hat{\mathbf{u}}}{n}-1\right)$: renormalization to stay away from the null vector

The eigenvalue is estimated as

$$\hat{\mu} = \frac{k_3}{k_2} \left(1 - \|\hat{\mathbf{u}}\|^2 \right)$$



Decentralized power iteration method (Yang et al. 2010; Zelazo et al. 2014)

$$\dot{\hat{\mathbf{u}}} = -k_1 T T^T \hat{\mathbf{u}} - k_2 M \hat{\mathbf{u}} - k_3 \left(\frac{\hat{\mathbf{u}}^T \hat{\mathbf{u}}}{n} - 1 \right)$$

connectivity maintenance

infinitesimal rigidity maintenance

The only remaining global quantities

- $T^T \hat{\mathbf{u}}$
- $\hat{\mathbf{u}}^T \hat{\mathbf{u}}$

can be estimated using the proportional/integral-average consensus estimator (PI-ACE) (Yang et al. 2010)



Possible limits of the gradient-based methods

- the robot could be unable to follow the gradient because of, e.g, input saturation
- possibility of local minima (depending on the environment complexity)



Possible **limits** of the decentralized methods:

- need for time-scale separation: decentralized estimator dynamics must be faster than motion control dynamics
- the gains of the decentralized estimator must be carefully tuned depending on N
- decentralized power iteration does not work for eigenvalues with multiplicity > 1
- (decentralized) power iteration has a relatively slow convergence

Possible destabilization due to non-perfect estimation can be mitigated using **passivity theory** (Robuffo Giordano et al. 2013)

Handling Multiple Objectives in Maintenance Problems



Connectivity in a network of robots is typically associated to

Inter-robot

- communication
- relative sensing

Quality of inter-robot sensing/communication modeled by a sufficiently smooth non-negative scalar function

$$\gamma_{ij} = \gamma(\mathbf{x}_i, \mathbf{x}_j, \mathbf{z}) \geq 0$$

Measures the quality of the mutual information exchange

- $\gamma_{ij} = 0$ if no exchange is possible and
- $\gamma_{ij} > 0$ otherwise
- the larger γ_{ij} the better the quality

Straightforward use:

$$\mathbf{w}_{ij} = \gamma_{ij}$$



In order to handle multiple objectives define

 $\mathbf{w}_{ij} = \alpha_{ij}\beta_{ij}\gamma_{ij}$

where

- $\alpha_{ij} \ge 0$ encodes hard constraints
- $\beta_{ij} \ge 0$ encodes soft requirements
- $\gamma_{ij} \ge 0$ encodes the communication/sensing objectives (defined before)

this defines the

- generalized connectivity, and a
- generalized infinitesimal rigidity



Hard constraints: conditions HD_1, HD_2, \ldots that must be true $\forall t \ge 0$

Maintenance methods automatically keep true a hard constraint: $HD_0 \equiv$ connectivity

Idea: define α_{ij} such that

• not HD_h for some $h \Rightarrow not HD_0$

How? Just define α_{ii} s.t.

• not HD_h for some $h \Rightarrow \alpha_{ij} = 0$, $\forall j = 1, \dots, N$

Why only $\alpha_{ij} = 0, \forall j = 1, \dots, N$?

- it is enough for non-connectivity $(\alpha_{ij} = 0, \forall j = 1, ..., N \text{ implies robot } i \text{ becomes disconnected from the rest})$
- is intrinsically decentralized

 α_{ij} must be smooth enough to allow for gradient computation

• the more $\alpha_{ij} \rightarrow 0$ the closer to not HD_h



Soft requirements: should be **preferably** realized by the individual pair (i, j)

Notes:

- gradient-based maintenance methods tend to maximize the maintenance eigenvalues (e.g., λ₂ or ς₇)
- maintenance eigenvalues monotonically increase w.r.t. $\mathbf{w}_{ij} \; orall (i,j) \in \mathcal{E}$

Idea: define β_{ij} such that

- has a unique maximum when the soft constraints are realized
- monotonically decreases down to $\beta_{ij} = 0$ otherwise

Non-perfect compliance with a soft requirement leads to

• corresponding decrease of maintenance eigenvalue

$$\downarrow \beta_{ij} \Rightarrow \downarrow \mathbf{w}_{ij} \Rightarrow \downarrow \lambda_2 \text{ (or } \downarrow \varsigma_7\text{)}$$

Complete violation of soft requirement

- leads to disconnected edge (i, j), but
- does not (in general) result in a global loss of connectivity for the graph

Applications

LAAS CNRS

Communication/sensing objectives $\rightarrow \gamma_{ij}(\mathbf{x}_i, \mathbf{x}_j, \mathbf{z})$

Proximity sensing model:

- D > 0 is a suitable sensing/communication maximum range (e.g., radio signal)
- robot *i* and *j* able to interact iff $\|\mathbf{x}_i \mathbf{x}_j\| < D$,

Proximity-visibility sensing model (e.g., onboard cameras):

- S_{ij} line-of-sight segment joining x_i and x_j
- robot *i* and *j* able to interact iff $\|\mathbf{x}_i \mathbf{x}_j\| < D$, and dist $(S_{ij(\mathbf{x}_i, \mathbf{x}_j)}, \text{obst}(\mathbf{z})) > D_{\text{vis}}$





Hard constraints $\rightarrow \alpha_{ij}$

e.g., inter-robot collision avoidance: $\|\mathbf{x}_i - \mathbf{x}_j\| > d_0$



Soft requirements $\rightarrow \beta_{ij}$

e.g., formation control, e.g., $\| \bm{x}_i - \bm{x}_j \| \simeq \textit{d}_{\rm des}$





Mission: **concurrent exploration** of a **sequence** of **targets** While maintaining "generalized" **connectivity**, i.e., including

- proximity/visibility sensing model
- collision avoidance
- preferred inter-distance

Connectivity maintenance in case of, e.g., second order systems:

$$\ddot{\mathbf{x}}_i = \left. \frac{\mathrm{d}V}{\mathrm{d}\mu} \right|_{\lambda_2(t)} \left. \frac{\partial\lambda_2}{\partial\mathbf{x}_i} \right|_{(\mathbf{x}_1,\dots,\mathbf{x}_N,\mathbf{z})} + u_i$$

$$u_i = -B\dot{\mathbf{x}}_i + f_i^{\mathrm{expl}}$$

- -Bx_i stabilizing damping
- f_i^{expl} multi-target exploration force (Nestmeyer et al. 2016)









Multi-Target Exploration with Connectivity





Multi-Target Exploration with Connectivity







Mission: **unilateral multi-user teleoperation** of some robots in the team While maintaining "generalized" **infinitesimal rigidity**, i.e., including

- proximity/visibility sensing model
- collision avoidance
- preferred inter-distance

Infinitesimal rigidity maintenance in case of, e.g., first order systems:

$$\begin{split} \dot{\mathbf{x}}_{i} &= \left. \frac{\mathrm{d}V}{\mathrm{d}\mu} \right|_{\varsigma_{7}(t)} \left. \frac{\partial\varsigma_{7}}{\partial \mathbf{x}_{i}} \right|_{(\mathbf{x}_{1},...,\mathbf{x}_{N},\mathbf{z})} + u_{i} \\ u_{i} &= \begin{cases} v_{i}^{h} & \text{if connected to a human} \\ 0 & \text{otherwise} \end{cases} \end{split}$$

• v_i^h desired velocity commanded by a human




Short summary

- Single scalars can define fundamental global properties
 - λ_2 Fiedler eigenvalue (Fiedler 1973)
 - ο _{\$7} rigidity eigenvalue (Zelazo et al. 2014)
- Distributed computation of the gradient is possible
 - + smooth
 - + online computation (fast)
 - presence of local minima

Some open problems

- coinciding eigenvalues
- local minima (using decentralized global planning?)



Bearing rigidity (from SE(3) to SE(3))

- D. Zelazo, Franchi, A., and Robuffo Giordano, P., "Rigidity Theory in SE(2) for Unscaled Relative Position Estimation using only Bearing", in 2014 European Control Conference, Strasbourg, France, 2014, pp. 2703-2708.
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- F. Schiano, Franchi A, Zelazo D, Robuffo Giordano P., "A Rigidity-Based Decentralized Bearing Formation Controller for Groups of Quadrotor UAVs", in 2016 IEEE/RSJ Int. Conf. on Intelligent Robots and System, Daejeon, South Korea, 2016.
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References



Questions?

Decentralized Estimation and Control Methods for Cooperative Robot Motion

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IEEE RAS Technical Committee on Multi-Robot Systems:

http://multirobotsystems.org/

- recently founded (Fall 2014)
- 330 members
- identifying and constantly tracking the common characteristics, problems, and achievements of multi-robot systems research in its several and diverse domains
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