# Trace polynomial optimization with applications in quantum information 

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Joint work with Felix Huber, Igor Klep and Jurij Volčič

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What is noncommutative optimization?

Eigenvalue optimization

Trace optimization

SDP hierarchies

Polynomial Bell inequalities

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Constraints

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\mathcal{D}_{S}=\left\{x: x_{i}^{2}=x_{i}, x_{i} x_{j}=x_{j} x_{i} \text { if } i \in\{1,2,3\}, j \in\{4,5,6\}\right\}
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\begin{aligned}
& \lambda_{\min }= \inf \left\{\langle f(x) \mathbf{v}, \mathbf{v}\rangle: x \in \mathcal{D}_{S},\|\mathbf{v}\|=1\right\} \\
&= \sup \lambda \\
& \text { s.t. } \quad f(x)-\lambda \mathbf{1} \succcurlyeq 0, \quad \forall x \in \mathcal{D}_{S}
\end{aligned}
$$

# Trace polynomials 

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f^{\star} & =x_{1}^{2} x_{2} x_{1}-\operatorname{tr}\left(x_{2}\right) \operatorname{tr}\left(x_{1} x_{2}\right) \operatorname{tr}\left(x_{1}^{2} x_{2}\right) x_{1} x_{2}
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$f^{\star} f$ hermitian square
$S \subset \operatorname{Sym} \mathbb{T} \quad x_{j}$ operators from finite von Neumann algebra
Constraints $\mathcal{D}_{S}=\{x: s(x) \succcurlyeq 0, \quad \forall s \in S\}$

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$\mathcal{A}=L^{\infty}(\mathcal{X}, \mu)$ is a vNa in $L^{2}(\mathcal{X}, \mu)$

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■ More complicated ones!

## Optimization over $\mathbb{T}$ : special cases

- Eigenvalue optimization no traces [Helton-McCallough 04, Navascuez-Pironio-Acin 08]

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- Univ case [Klep-Pascoe-Volcic 21]: $f \succcurlyeq 0 \Rightarrow f=$ SOHS/SOHS
- Multilinear case [Huber 21]


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classical correlations = convex combinations of deterministic correlations

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Clauser-Horne-Shimony-Holt (CHSH) inequality is violated by quantum systems:
$-P_{a}(1 \mid 0)-P_{b}(1 \mid 0)+P(1,1 \mid 0,0)+P(1,1 \mid 0,1)+P(1,1 \mid 1,0)-P(1,1 \mid 1,1) \leqslant 0$

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Alice \& Bob share a bipartite quantum state $\Psi$ and they answer $s, t$ by performing quantum measurements on their part of $\Psi$ :

$$
P(a, b \mid s, t)=\Psi^{\star} X_{s}^{a} Y_{t}^{b} \Psi, \quad P(a \mid s)=\Psi^{\star} X_{s}^{a} \Psi, \quad P(b \mid t)=\Psi^{\star} Y_{t}^{b} \Psi
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Clauser-Horne-Shimony-Holt (CHSH) inequality is violated by quantum systems:
$-P_{a}(1 \mid 0)-P_{b}(1 \mid 0)+P(1,1 \mid 0,0)+P(1,1 \mid 0,1)+P(1,1 \mid 1,0)-P(1,1 \mid 1,1) \leqslant 0$
Alice \& Bob share a bipartite quantum state $\Psi$ and they answer $s, t$ by performing quantum measurements on their part of $\Psi$ :

$$
P(a, b \mid s, t)=\Psi^{\star} X_{s}^{a} Y_{t}^{b} \Psi, \quad P(a \mid s)=\Psi^{\star} X_{s}^{a} \Psi, \quad P(b \mid t)=\Psi^{\star} Y_{t}^{b} \Psi
$$

$X_{s}^{a}, Y_{t}^{b}$ are bounded operators on separable Hilbert spaces s.t.:

$$
\begin{array}{rll}
X_{s}^{a} Y_{t}^{b}=Y_{t}^{b} X_{s}^{a}, & X_{s}^{a} X_{s}^{a}=X_{s}^{a}, & Y_{t}^{b} Y_{t}^{b}=Y_{t}^{b} \\
X_{s}^{a} X_{s}^{a^{\prime}}=Y_{t}^{b} Y_{t}^{b^{\prime}}=0, & \sum_{a} X_{s}^{a}=\sum_{b} Y_{t}^{b}=I &
\end{array}
$$

"Mathematically": inequality on eigenvalues of noncommutative polynomials

## Motivation: Bell inequalities

Entanglement in quantum mechanics
$\rightarrow$ upper bounds for violation levels of Bell inequalities

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Entanglement in quantum mechanics
$\rightarrow$ upper bounds for violation levels of Bell inequalities
[Pozas et al 19] extension $\rightarrow$ identify correlations not attainable in entanglement-swapping scenario (quantum teleportation)

Quantum physics operators $x_{i}, y_{j}$ satisfy causal constraints:

$$
\operatorname{tr}\left(x_{1} x_{2} y_{1} y_{2}\right)-\operatorname{tr}\left(x_{1} x_{2}\right) \operatorname{tr}\left(y_{1} y_{2}\right)=0 .
$$

## What is noncommutative optimization?

Eigenvalue optimization

## Trace optimization

## SDP hierarchies

Polynomial Bell inequalities

## Eigenvalue optimization

$$
\lambda_{\min }=\inf _{\mathbf{v}, \mathcal{H}}\left\{\langle f(x) \mathbf{v}, \mathbf{v}\rangle: x \in \mathcal{D}_{S},\|\mathbf{v}\|=1, \mathbf{v} \in \mathcal{H}\right\}
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&= \sup \quad \lambda \\
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$\mathcal{M}(S)$ Archimedean quadratic module: $N-\sum_{i} x_{i}^{2} \succcurlyeq 0$
Theorem: NC Putinar's representation [Helton-McCullough 02]
$f \succcurlyeq 0$ on $\mathcal{D}_{S} \Longrightarrow f+\varepsilon \in \mathcal{M}(S)$, for all $\varepsilon>0$

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$f-\lambda \mathbf{1}=\sum_{i} h_{i}^{\star} h_{i}+\sum_{s} \sum_{i} t_{s i}^{\star} s t_{s i}$ with $h_{i}$, $t_{s i}$ of bounded degrees

## Sparse eigenvalue optimization

Self-adjoint noncommutative (NC) variables $x=\left(x_{1}, \ldots, x_{n}\right)$
Theorem [Helton \& McCullough '02]
$f \succcurlyeq 0 \Leftrightarrow f$ SOS (all positive polynomials are sums of squares)

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## Sparse eigenvalue optimization

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## Theorem: Sparse NC Positivstellensatz [Klep Magron Povh '21]

$f=\sum_{k} f_{k}, f_{k}$ depends on $x\left(I_{k}\right)$
$f>0$ on $\mathcal{D}_{S}$
Each $g_{j}$ depends on some $I_{k}$
RIP holds for $\left(I_{k}\right)$
Ball constraints for each $x\left(I_{k}\right)$

$$
f=\sum_{k, i}\left(s_{k i}^{\star} s_{k i}+\sum_{j \in J_{k}} t_{j i}{ }^{\star} g_{j} t_{j i}\right)
$$

$s_{k i}$ "sees" vars in $I_{k}$
$t_{j i}$ "sees" vars from $g_{j}$

## Sparse eigenvalue optimization

$\mathbf{I}_{3322}$ Bell inequality (entanglement in quantum information)

$$
f=x_{1}\left(x_{4}+x_{5}+x_{6}\right)+x_{2}\left(x_{4}+x_{5}-x_{6}\right)+x_{3}\left(x_{4}-x_{5}\right)-x_{1}-2 x_{4}-x_{5}
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Maximal violation levels $\rightarrow$ upper bounds on $\lambda_{\text {max }}$ of $f$ on $\left\{x: x_{i}^{2}=x_{i}, x_{i} x_{j}=x_{j} x_{i}\right.$ if $\left.i \in\{1,2,3\}, j \in\{4,5,6\}\right\}$

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& 20.2550008 \\
& 0.2509397
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3'
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$6 \quad 0.2508753977180$
(1 hour)

Performance


Accuracy

## What is noncommutative optimization?

## Eigenvalue optimization

Trace optimization

## SDP hierarchies

Polynomial Bell inequalities

## Trace optimization

$$
\begin{aligned}
\operatorname{tr}_{\min }= & \inf \left\{\operatorname{tr}(f(x)): x \in D_{S}\right\} \\
= & \sup \quad m \\
& \text { s.t. } \quad \operatorname{tr}(f(x)-m) \geqslant 0, \quad \forall x \in \mathcal{D}_{S}
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$\operatorname{tr}_{\text {min }}^{\mathrm{II}_{1}}=$ minimal trace over the union of type $-\mathrm{II}_{1} \mathrm{vN}$ algebras
". Disproving Connes' embedding conjecture: $\operatorname{tr}_{\min }^{\mathrm{II}_{1}}<\operatorname{tr}_{\text {min }}$

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Converging hierarchy with cyclic quadratic modules: "̈ replace " $\operatorname{tr}(f-m) \geqslant 0$ on $\mathcal{D}_{S}^{\mathrm{II}_{1} "}$ by $f-m \mathbf{1} \in \mathcal{M}^{\text {cyc }}(S)_{r}$
$\mathcal{M}^{\text {cyc }}(S)_{r}=$ polynomials with same trace as some from $\mathcal{M}(S)_{r}$

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$\mathcal{M}^{\text {cyc }}(S)_{r}=$ polynomials with same trace as some from $\mathcal{M}(S)_{r}$ How to extend it to sums of trace products $T$ ?

## Kadison-Dubois representation theorem

$\chi_{\mathcal{M}}:=\left\{\varphi: \mathbf{T} \rightarrow \mathbb{R} \mid \varphi\right.$ homomorphism, $\left.\varphi(\mathcal{M}) \subseteq \mathbb{R}_{\geqslant 0,} \varphi(1)=1\right\}$

## Kadison-Dubois representation theorem

$\chi_{\mathcal{M}}:=\left\{\varphi: \mathbf{T} \rightarrow \mathbb{R} \mid \varphi\right.$ homomorphism, $\left.\varphi(\mathcal{M}) \subseteq \mathbb{R}_{\geqslant 0,} \varphi(1)=1\right\}$
Theorem: Kadison-Dubois [Marshall 08]
Given an Archimedean quadratic module $\mathcal{M} \subseteq \mathbf{T} \& f \in \mathrm{~T}$ :

$$
\forall \varphi \in \chi_{\mathcal{M}} \quad \varphi(f) \geqslant 0 \quad \Leftrightarrow \quad \forall \varepsilon>0 \quad f+\varepsilon \in \mathcal{M}
$$

## Representation of positive elements of $T$

For $S \subseteq \mathrm{~T}$, "augment" $S$ with traces of hermitian squares:

$$
S(N)=S \cup\left\{\operatorname{tr}\left(p p^{\star}\right) \mid p \in \mathbb{R}\langle\underline{x}\rangle\right\} \cup\left\{N^{k}-\operatorname{tr}\left(x_{j}^{2 k}\right) \mid k \in \mathbb{N}\right\} \subseteq \mathrm{T}
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Elements of $\mathcal{M}(S(N))$ are

$$
p_{1}^{2} s \quad p_{2}^{2}\left(N^{k}-\operatorname{tr}\left(x_{j}^{2 k}\right)\right) \quad \operatorname{tr}\left(a a^{\star}\right)
$$

for $s \in S, p_{i} \in \mathbf{T}, a \in \mathbb{T}$

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Lemma [Klep-M.-Volcic 20]
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for $s \in S, p_{i} \in \mathrm{~T}, a \in \mathbb{T}$
Lemma [Klep-M.-Volcic 20]
$\mathcal{M}(S(N))$ is archimedean
Proof
By induction: $\forall$ word $w, m \pm \operatorname{tr}(w) \in \mathcal{M}(S(N))$ for some $m>0$

$$
w=x_{j}^{2 k} \Longrightarrow N^{k}+1+2 \operatorname{tr}\left(x_{j}^{k}\right)=\left(N^{k}-\operatorname{tr}\left(x_{j}^{2 k}\right)\right)+\operatorname{tr}\left(\left(x_{j}^{k}+1\right)^{2}\right)
$$

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\begin{gathered}
S(N)=S \cup\left\{\operatorname{tr}\left(p p^{\star}\right) \mid p \in \mathbb{R}(\underline{x})\right\} \cup\left\{N^{k}-\operatorname{tr}\left(x_{j}^{2 k}\right) \mid k \in \mathbb{N}\right\} \subseteq T \\
S[N]=S \cup\left\{N-x_{j}^{2}\right\} \subset \mathbb{T}
\end{gathered}
$$

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Theorem [Klep-M.-Volcic 20]
$f \geqslant 0$ on $\mathcal{D}_{S[N]}^{\mathrm{I}_{1}} \quad \Leftrightarrow \quad f+\varepsilon \in \mathcal{M}(S(N))$ for all $\varepsilon>0$

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## Tracial words \& moment matrices

$\mathbb{T}$-words $=\left\{\prod_{i} \operatorname{tr}\left(u_{i}\right) v: u_{i}, v\right.$ words $\}$ and T-words $\operatorname{tr}\left(x_{1}\right)^{2}$ is a T -word, $\operatorname{tr}\left(x_{1}\right) x_{1}$ is a $\mathbb{T}$-word

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- Tracial degree $=$ up to cyclic equivalence
$\mathbf{W}_{r}^{\mathbb{T}}=$ vector of $\mathbb{T}$-words of with tracial degree $\leqslant r$
$n=1$ : $\mathbf{W}_{2}^{\mathbb{T}}$ contains $1, x_{1}, x_{1}^{2}, \operatorname{tr}\left(x_{1}\right), \operatorname{tr}\left(x_{1}^{2}\right), \operatorname{tr}\left(x_{1}\right) x_{1}$


## Tracial words \& moment matrices

$\mathbb{T}$-words $=\left\{\prod_{i} \operatorname{tr}\left(u_{i}\right) v: u_{i}, v\right.$ words $\}$ and T-words $\operatorname{tr}\left(x_{1}\right)^{2}$ is a $\mathbb{T}$-word, $\operatorname{tr}\left(x_{1}\right) x_{1}$ is a $\mathbb{T}$-word
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## SDP hierarchy for T

## Reminder:

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S(N)=S \cup\left\{\operatorname{tr}\left(p p^{\star}\right) \mid p \in \mathbb{R}\langle\underline{x}\rangle\right\} \cup\left\{N^{k}-\operatorname{tr}\left(x_{j}^{2 k}\right) \mid k \in \mathbb{N}\right\} \subseteq \mathrm{T}
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\mathcal{M}(S(N))_{r}=\left\{\sum_{i} p_{i}^{2} s_{i}: s_{i} \in S(N), p_{i} \in \mathrm{~T}, \operatorname{deg}\left(p_{i}^{2} s_{i}\right) \leq 2 r\right\}
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Elements of $\mathcal{M}(S(N))_{r}$ are

$$
p_{1}^{2} s \quad p_{2}^{2}\left(N^{k}-\operatorname{tr}\left(x_{j}^{2 k}\right)\right) \quad \operatorname{tr}\left(a a^{\star}\right)
$$

for $s \in S, p_{i} \in \mathrm{~T}, a \in \mathbb{T}$

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## Theorem [Klep-M.-Volcic 20]

There is no duality gap and $f_{r} \rightarrow f_{\text {min }}^{\mathrm{II}_{1}}$ as $r \rightarrow \infty$

## SDP hierarchy for $\mathbb{T}$

$S \subset$ Sym $\mathbb{T}$
$\ddot{\nabla}$ Reduction from the general trace setting to the pure trace

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## What is noncommutative optimization?

Eigenvalue optimization

Trace optimization

SDP hierarchies

Polynomial Bell inequalities

## Polynomial Bell inequalities

## Classical world

$$
\psi^{*}\left(A_{1} \otimes B_{1}+A_{1} \otimes B_{2}+A_{2} \otimes B_{1}-A_{2} \otimes B_{2}\right) \psi \leqslant 2
$$

for separable states $\psi \in \mathbb{R}^{k} \otimes \mathbb{R}^{k}$ and matrices $A_{j}, B_{j}$ satisfying

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for separable states but ... 5 for one maximally entangled state

- 2nd dense SDP relaxation of the corresponding trace problem outputs $5=$ max value for all maximal entangled states
- 2nd sparse SDP gives also $5 \ldots 10$ times faster


## Conclusion and perspectives

CONVERGING HIERARCHIES to minimize pure trace polynomials

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## Thank you for your attention!

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