Trace polynomial optimization with applications in quantum information

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Joint work with Felix Huber, Igor Klep and Jurij Volčič

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Eigenvalue optimization

Trace optimization

SDP hierarchies

Polynomial Bell inequalities

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Constraints

 $\mathcal{D}_{S} = \{x : x_{i}^{2} = x_{i}, x_{i}x_{j} = x_{j}x_{i} \text{ if } i \in \{1, 2, 3\}, j \in \{4, 5, 6\}\}$

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Trace polynomials

Elements of $\mathbb{T} = \mathsf{T}\langle x \rangle$

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 $f^{\star}f$ hermitian square

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 f^*f hermitian square $S \subset \text{Sym } \mathbb{T}$ x_j operators from finite von Neumann algebra Constraints $\mathcal{D}_S = \{x : s(x) \geq 0, \forall s \in S\}$

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Commutant A' = elements commuting with elements of A

If $\mathcal{A} = \mathcal{B}(\mathcal{H})$ then $\mathcal{A}' = \mathbb{C}Id_{\mathcal{H}}$ if "Trivial commutant" = factor

$$\mathcal{A} = L^{\infty}(\mathcal{X}, \mu)$$
 is a vNa in $L^{2}(\mathcal{X}, \mu)$

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- More complicated ones!

Optimization over T: **special cases**

■ Eigenvalue optimization V no traces [Helton-McCallough 04, Navascuez-Pironio-Acin 08]

$$\begin{split} \lambda_{\min} &= \inf \left\{ \langle f(x) \mathbf{v}, \mathbf{v} \rangle : x \in \mathcal{D}_{S}, \|\mathbf{v}\| = 1 \right\} \\ &= \sup \left\{ \lambda : f(x) - \lambda \mathbf{1} \succcurlyeq \mathbf{0}, \quad \forall x \in \mathcal{D}_{S} \right\} \end{split}$$

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- Univ case [Klep-Pascoe-Volcic 21]: $f \succcurlyeq 0 \Rightarrow f = \text{SOHS/SOHS}$
- Multilinear case [Huber 21]

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Trace polynomial optimization with applications in quantum information

"Pillars" of quantum physics: violations imply that properties (e.g. entanglement) can't be represented by classical physics "Pillars" of quantum physics: violations imply that properties (e.g. entanglement) can't be represented by classical physics

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classical correlations = convex combinations of deterministic correlations

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Clauser-Horne-Shimony-Holt (CHSH) inequality is violated by quantum systems:

 $-P_a(1|0) - P_b(1|0) + P(1,1|0,0) + P(1,1|0,1) + P(1,1|1,0) - P(1,1|1,1) \leq 0$

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Alice & Bob share a bipartite quantum state Ψ and they answer *s*, *t* by performing quantum measurements on their part of Ψ :

 $P(a,b|s,t) = \Psi^{\star} X_s^a Y_t^b \Psi, \quad P(a|s) = \Psi^{\star} X_s^a \Psi, \quad P(b|t) = \Psi^{\star} Y_t^b \Psi$

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 X_s^a , Y_t^b are bounded operators on separable Hilbert spaces s.t.:

$$\begin{aligned} X_{s}^{a}Y_{t}^{b} &= Y_{t}^{b}X_{s}^{a}, \quad X_{s}^{a}X_{s}^{a} = X_{s}^{a}, \\ X_{s}^{a}X_{s}^{a'} &= Y_{t}^{b}Y_{t}^{b'} = 0, \quad \sum_{a}X_{s}^{a} = \sum_{b}Y_{t}^{b} = I \end{aligned}$$

"Mathematically": inequality on eigenvalues of noncommutative

polynomials

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Entanglement in quantum mechanics

 \rightarrow upper bounds for violation levels of Bell inequalities

Entanglement in quantum mechanics

 \rightarrow **upper bounds** for violation levels of Bell inequalities

[Pozas et al 19] extension \rightarrow identify correlations not attainable in entanglement-swapping scenario (quantum teleportation)

Quantum physics operators x_i , y_j satisfy causal constraints:

$$\operatorname{tr}(x_1x_2y_1y_2) - \operatorname{tr}(x_1x_2)\operatorname{tr}(y_1y_2) = 0.$$

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What is noncommutative optimization?

Eigenvalue optimization

Trace optimization

SDP hierarchies

Polynomial Bell inequalities

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 $\mathcal{M}(S)$ Archimedean quadratic module: $N - \sum_i x_i^2 \geq 0$

Theorem: NC Putinar's representation [Helton-McCullough 02] $f \succeq 0 \text{ on } \mathcal{D}_S \implies f + \varepsilon \in \mathcal{M}(S)$, for all $\varepsilon > 0$

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NC variant of Lasserre's Hierarchy for λ_{\min} : \forall replace " $f - \lambda \mathbf{1} \succeq 0$ on \mathcal{D}_{S} " by $f - \lambda \mathbf{1} \in \mathcal{M}(S)_{r}$

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$$= \sup \quad \lambda$$
$$s.t. \quad f(x) - \lambda \mathbf{1} \succeq 0, \quad \forall x \in \mathcal{D}_S$$

 $\mathcal{M}(S)$ Archimedean quadratic module: $N - \sum_i x_i^2 \succcurlyeq 0$

Theorem: NC Putinar's representation [Helton-McCullough 02] $f \succeq 0 \text{ on } \mathcal{D}_S \implies f + \varepsilon \in \mathcal{M}(S)$, for all $\varepsilon > 0$

NC variant of Lasserre's Hierarchy for λ_{\min} : $\forall replace "f - \lambda \mathbf{1} \geq 0 \text{ on } \mathcal{D}_S"$ by $f - \lambda \mathbf{1} \in \mathcal{M}(S)_r$ $f - \lambda \mathbf{1} = \sum_i h_i^* h_i + \sum_s \sum_i t_{si}^* st_{si}$ with h_i , t_{si} of **bounded** degrees

Self-adjoint noncommutative (NC) variables $x = (x_1, ..., x_n)$

Theorem [Helton & McCullough '02]

 $f \succcurlyeq 0 \Leftrightarrow f$ SOS (all positive polynomials are sums of squares)

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Theorem: Sparse NC Positivstellensatz [Klep Magron Povh '21]

$$f = \sum_{k} f_{k}, f_{k} \text{ depends on } x(I_{k})$$

$$f > 0 \text{ on } \mathcal{D}_{S}$$

Each g_{j} depends on some I_{k}
RIP holds for $(I_{k}) \implies$
Ball constraints for each $x(I_{k})$

$$f = \sum_{k,i} (s_{ki}^{\star} s_{ki} + \sum_{j \in J_k} t_{ji}^{\star} g_j t_{ji})$$

$$s_{ki} \text{ "sees" vars in } I_k$$

$$t_{ji} \text{ "sees" vars from } g_j$$

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Trace polynomial optimization with applications in quantum information

I₃₃₂₂ Bell inequality (entanglement in quantum information)

 $f = x_1(x_4 + x_5 + x_6) + x_2(x_4 + x_5 - x_6) + x_3(x_4 - x_5) - x_1 - 2x_4 - x_5$

Maximal violation levels \rightarrow **upper bounds** on λ_{\max} of f on $\{x : x_i^2 = x_i, x_i x_j = x_j x_i \text{ if } i \in \{1, 2, 3\}, j \in \{4, 5, 6\}\}$

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level	sparse
2	0.2550008

dense [Pál & Vértesi '18] 0.2509397

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level	sparse	dense [Pál & Vértesi
2	0.2550008	0.2509397
3	0.2511592	0.2508756
3'		0.25087 <mark>54</mark> (1 day)
4	0.2508917	

5 0.25087<mark>63</mark>

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Trace polynomial optimization with applications in quantum information

What is noncommutative optimization?

Eigenvalue optimization

Trace optimization

SDP hierarchies

Polynomial Bell inequalities

Trace optimization

$$tr_{\min} = \inf\{tr(f(x)) : x \in D_S\}$$

= sup m
s.t. $tr(f(x) - m) \ge 0$, $\forall x \in \mathcal{D}_S$

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$$\label{eq:trmin} \begin{split} tr_{min}^{II_1} = \mbox{minimal trace over the union of type} - II_1 \ \mbox{vN algebras} \\ \overleftarrow{\mbox{v}} \ \mbox{Disproving Connes' embedding conjecture: } tr_{min}^{II_1} < tr_{min} \end{split}$$
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Converging hierarchy with cyclic quadratic modules: \forall replace "tr $(f - m) \ge 0$ on $\mathcal{D}_{S}^{\Pi_{1}}$ " by $f - m\mathbf{1} \in \mathcal{M}^{\text{cyc}}(S)_{r}$

 $\mathcal{M}^{cyc}(S)_r$ = polynomials with same trace as some from $\mathcal{M}(S)_r$

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$\mathcal{M}^{cyc}(S)_r$ = polynomials with same trace as some from $\mathcal{M}(S)_r$ How to extend it to sums of trace products T?

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Kadison-Dubois representation theorem

 $\chi_{\mathcal{M}} := \{ \varphi : \mathsf{T} \to \mathbb{R} \mid \varphi \text{ homomorphism}, \varphi(\mathcal{M}) \subseteq \mathbb{R}_{\geqslant 0}, \ \varphi(1) = 1 \}$

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Theorem: Kadison-Dubois [Marshall 08]

Given an Archimedean quadratic module $\mathcal{M} \subseteq \mathsf{T} \& f \in \mathsf{T}$:

$$\forall \varphi \in \chi_{\mathcal{M}} \quad \varphi(f) \geqslant 0 \qquad \Leftrightarrow \qquad \forall \varepsilon > 0 \quad f + \varepsilon \in \mathcal{M}$$

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Trace polynomial optimization with applications in quantum information

For $S \subseteq T$, "augment" S with traces of hermitian squares:

 $S(N) = S \cup \{ \operatorname{tr}(pp^{\star}) \mid p \in \mathbb{R} \langle \underline{x} \rangle \} \cup \{ N^k - \operatorname{tr}(x_j^{2k}) \mid k \in \mathbb{N} \} \subseteq \mathsf{T}$

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Elements of $\mathcal{M}(S(N))$ are

$$p_1^2 s \quad p_2^2 \left(N^k - \operatorname{tr}(x_i^{2k}) \right) \quad \operatorname{tr}(aa^{\star})$$

for $s \in \mathbf{S}$, $p_i \in \mathsf{T}$, $a \in \mathbb{T}$

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Lemma [Klep-M.-Volcic 20]

 $\mathcal{M}(S(N))$ is archimedean

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Lemma [Klep-M.-Volcic 20]

 $\mathcal{M}(S(N))$ is archimedean

Proof

By induction: \forall word w, $m \pm tr(w) \in \mathcal{M}(S(N))$ for some m > 0

$$w = x_j^{2k} \implies N^k + 1 + 2\operatorname{tr}(x_j^k) = (N^k - \operatorname{tr}(x_j^{2k})) + \operatorname{tr}((x_j^k + 1)^2)$$

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Theorem [Klep-M.-Volcic 20] $f \ge 0$ on $\mathcal{D}_{S[N]}^{II_1}$ \Leftrightarrow $f + \varepsilon \in \mathcal{M}(S(N))$ for all $\varepsilon > 0$

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Trace polynomial optimization with applications in quantum information

What is noncommutative optimization?

Eigenvalue optimization

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SDP hierarchies

Polynomial Bell inequalities

 \mathbb{T} -words = { $\prod_i \operatorname{tr}(u_i)v : u_i, v \text{ words}$ } and T-words tr $(x_1)^2$ is a T-word, tr $(x_1)x_1$ is a \mathbb{T} -word

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 $\mathbf{W}_r^{\mathbb{T}}$ = vector of \mathbb{T} -words of with tracial degree $\leq r$ n = 1: $\mathbf{W}_2^{\mathbb{T}}$ contains 1, $x_1, x_1^2, \operatorname{tr}(x_1), \operatorname{tr}(x_1^2), \operatorname{tr}(x_1)x_1$

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Tracial moment matrix $\mathbf{M}_r^{\mathbb{T}}(L)$ for a linear functional $L: \mathsf{T} \to \mathbb{R}$:

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Trace localizing matrix $\mathbf{M}_r^{\mathbb{T}}(sL)$ for $s \in \mathbb{T}$:

• indexed by $\mathbf{W}_r^{\mathbb{T}}$

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Trace polynomial optimization with applications in quantum information

Reminder:

$$S(N) = S \cup \{ \operatorname{tr}(pp^{\star}) \mid p \in \mathbb{R} \langle \underline{x} \rangle \} \cup \{ N^{k} - \operatorname{tr}(x_{j}^{2k}) \mid k \in \mathbb{N} \} \subseteq \mathsf{T}$$

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$$\begin{split} & \inf_{\text{linear } L} \quad L(f) \\ & \text{s.t.} \quad (\mathbf{M}_r^{\mathbb{T}}(L))_{u,v} = (\mathbf{M}_r^{\mathbb{T}}(L))_{w,z} \quad \text{whenever } \operatorname{tr}(u^\star v) = \operatorname{tr}(w^\star z) \\ & \quad (\mathbf{M}_r^{\mathbb{T}}(L))_{1,1} = 1 \\ & \quad \mathbf{M}_r^{\mathbb{T}}(L) \succcurlyeq 0, \\ & \quad \mathbf{M}_{r-r_s}^{\mathbb{T}}(s\,L) \succcurlyeq 0, \quad \text{ for all } s \in S \\ & \quad \mathbf{M}_{r-k}^{\mathbb{T}}((N^k - \operatorname{tr}(x_i^{2k}))\,L) \succcurlyeq 0 \end{split}$$

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Theorem [Klep-M.-Volcic 20]

There is no duality gap and $f_r \rightarrow f_{\min}^{\text{II}_1}$ as $r \rightarrow \infty$

SDP hierarchy for $\mathbb T$

 $S \subset \text{Sym } \mathbb{T}$ \bigvee Reduction from the general trace setting to the pure trace

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$$\psi^*(A_1 \otimes B_1 + A_1 \otimes B_2 + A_2 \otimes B_1 - A_2 \otimes B_2)\psi \leqslant 2$$

for separable states $\psi \in \mathbb{R}^k \otimes \mathbb{R}^k$ and matrices A_j , B_j satisfying $A_j^* = A_j$, $A_j^2 = I$, $B_j^* = B_j$, $B_j^2 = I$

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TSIRELSON'S BOUND for maximally entangled states $\psi = \frac{1}{\sqrt{k}} \sum_{j=1}^{k} e_j \otimes e_j \in \mathbb{R}^k \otimes \mathbb{R}^k$

$$\psi^*(A_1 \otimes B_1 + A_1 \otimes B_2 + A_2 \otimes B_1 - A_2 \otimes B_2)\psi \leqslant 2$$

for separable states $\psi \in \mathbb{R}^k \otimes \mathbb{R}^k$ and matrices A_j , B_j satisfying $A_j^* = A_j$, $A_j^2 = I$, $B_j^* = B_j$, $B_j^2 = I$

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$$2 \to 2\sqrt{2} = \operatorname{tr}_{\max}\{a_1b_1 + a_1b_2 + a_2b_1 - a_2b_2 : a_j^2 = b_j^2 = 1\}$$

Victor Magron

Trace polynomial optimization with applications in quantum information

Polynomial Bell inequalities

COVARIANCES OF QUANTUM CORRELATIONS

$$\operatorname{cov}_{\psi}(X,Y) := \psi^*(X \otimes Y)\psi - \psi^*(X \otimes I)\psi \cdot \psi^*(I \otimes Y)\psi$$

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for separable states but ...

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V 2nd dense SDP relaxation of the corresponding trace problem outputs 5

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for separable states but ... 5 for one maximally entangled state

 \checkmark 2nd dense SDP relaxation of the corresponding trace problem outputs $5 = \max$ value for **all** maximal entangled states

V 2nd sparse SDP gives also 5 ... 10 times faster

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Trace polynomial optimization with applications in quantum information

Implementation available in github:TSSOS

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Hierarchy for minimal eigenvalue problem, degree bounds?

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Thank you for your attention!

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