

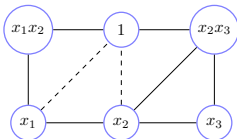
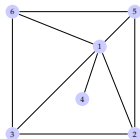
# Exploiting Correlative and Term Sparsity in Noncommutative Polynomial Optimization

Victor Magron (LAAS CNRS)

Joint work with Igor Klep, Janez Povh and Jie Wang

SIAM Conference on Optimization

23 July 2021



# What is a noncommutative polynomial?

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An element of  $\mathbb{R}\langle \underline{x} \rangle$

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Symmetric noncommutative variables  $\underline{x} = (x_1, \dots, x_n)$

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with  $x_1 x_2 \neq x_2 x_1$ , **involution**  $(x_1 x_2)^* = x_2 x_1$

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$S \subset \text{Sym } \mathbb{R}\langle \underline{x} \rangle$   $X_j \in$  finite von Neumann algebra

Constraints:  $\mathcal{D}_S = \{ \underline{X} = (X_1, \dots, X_n) : \mathbf{s}(\underline{X}) \succcurlyeq 0, \quad \forall \mathbf{s} \in S \}$

# Motivation: quantum information

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Entanglement in quantum mechanics

→ **upper bounds** for violation levels of Bell inequalities

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Finding violation → solving a **maximal** eigenvalue problem!

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[Werner '89] Werner witnesses to separate pure/entangled states

# Motivation: condensed matter

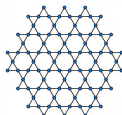
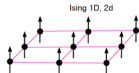
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Ground-state energy  $\Leftrightarrow$  minimal eigenvalue of an Hamiltonian

$$f = \sum_{\langle i,j \rangle} (x_i x_j + y_i y_j + z_i z_j)$$

spin states  $(x_i, y_i, z_i)$ , constraints

Lattices: 1D 2D Kagome



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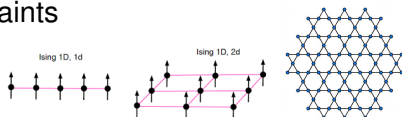
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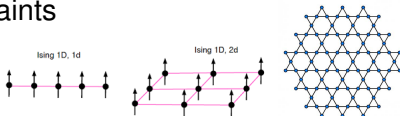
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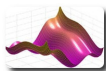
Existing  $\pm$  efficient techniques: quantum Monte Carlo & variational algorithms  $\Rightarrow$  **upper bounds** on minimal energy

# Fighting the scalability issue

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NP-hard NON CONVEX Problem  $f_{\min} = \inf f(x)$

## Theory



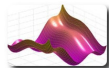
INFINITE LP

$$\begin{aligned} & \sup \lambda \\ \Leftrightarrow & \text{with } f - \lambda \geq 0 \end{aligned}$$

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## Practice



with  $f - \lambda = \text{sum of squares}$

**FINITE SDP**

$\Leftarrow$  **fixed** degree

HIERARCHY of **CONVEX** (semidefinite) **PROBLEMS**  $\uparrow f_{\min}$

[Lasserre '01]

[Navascués Pironio Acín '08] for NC problems

✓ Relaxations  $\implies$  **certified lower bounds**

✗ degree  $d$  &  $n$  vars  $\implies n^{2d}$  variables



Moment-sums of squares hierarchies

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Semialgebraic set  $\mathcal{D}_S = \{\mathbf{x} \in \mathbb{R}^n : s_j(\mathbf{x}) \geq 0\}$

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$$\text{Quadratic module: } \mathcal{M}(S)_d = \left\{ \sigma_0 + \sum_j \sigma_j s_j, \text{ deg } \sigma_j s_j \leq 2d \right\}$$

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Can be computed with SDP solvers (CSDP, SDPA, MOSEK)

**“No Free Lunch” Rule:**  $\binom{n+2r}{n}$  SDP variables

# Eigenvalue optimization

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$\mathcal{M}(S)$  Archimedean quadratic module:  $N - \sum_i x_i^2 \succcurlyeq 0$

Theorem: NC Putinar's Psatz [Helton-McCullough 02]

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$f - \lambda \mathbf{1} = \sum_i h_i^* h_i + \sum_j \sum_i t_{ji}^* s_j t_{ji}$  with  $h_i, t_{ji}$  of **bounded** degrees

# Trace optimization

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$$\text{tr}_{\min} = \inf\{\text{tr}(f(\underline{X})) : \underline{X} \in \mathcal{D}_S\}$$

$$= \sup m$$

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$\mathrm{tr}_{\min}^{\mathrm{II}_1}$  = minimal trace over the union of type  $\mathrm{II}_1$  vN algebras

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Converging hierarchy with cyclic quadratic modules:

💡 replace " $\mathrm{tr}(f - m) \geq 0$  on  $\mathcal{D}_S^{\mathrm{II}_1}$ " by  $f - m\mathbf{1} \in \mathcal{M}^{\mathrm{cyc}}(S)_d$

$\mathcal{M}^{\mathrm{cyc}}(S)_d$  = polynomials with same trace as some from  $\mathcal{M}(S)_d$

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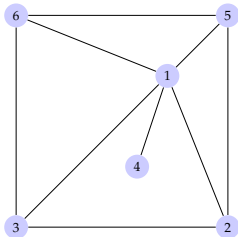
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💡 Exploit few links between **variables** [Lasserre, Waki et al. '06]

$$x_2x_5 + x_3x_6 - x_2x_3 - x_5x_6 + x_1(-x_1 + x_2 + x_3 - x_4 + x_5 + x_6)$$

Chordal graph  $G$

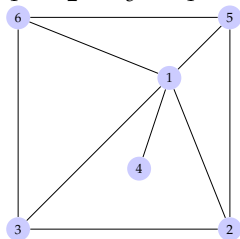


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maximal cliques  $I_k$

$$I_1 = \{1, 4\}$$

$$I_2 = \{1, 2, 3, 5\}$$

$$I_3 = \{1, 3, 5, 6\}$$

Dense SDP: 210 vars

Sparse SDP: 115 vars

Average size  $\kappa \rightsquigarrow \kappa^{2d}$  vars

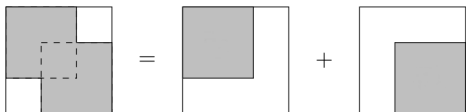
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## Theorem [Griewank Toint '84]

Chordal graph  $G$  with maximal cliques  $I_1, I_2$

$Q_G \succcurlyeq 0$  with nonzero entries at edges of  $G$

$\implies Q_G = P_1^T Q_1 P_1 + P_2^T Q_2 P_2$  with  $Q_k \succcurlyeq 0$  indexed by  $I_k$



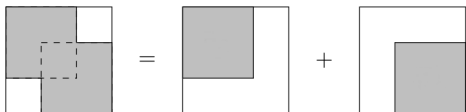
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Sparse  $f = f_1 + f_2$  where  $f_k$  involves **only** variables in  $I_k$

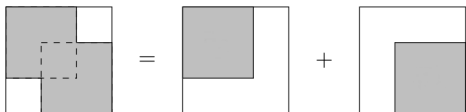
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## Theorem: Sparse Putinar's representation [Lasserre '06]

$f > 0$  on  $\{x : s_j(x) \geq 0\}$

chordal graph  $G$  with cliques  $I_k \implies$

ball constraints for each  $x(I_k)$

$$f = \sigma_{01} + \sigma_{02} + \sum_j \sigma_j s_j$$

SOS  $\sigma_{0k}$  "sees" vars in  $I_k$

$\sigma_j$  "sees" vars from  $s_j$

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symmetric noncommutative (NC) variables  $\underline{x} = (x_1, \dots, x_n)$

**Theorem [Helton-McCullough 02]**

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**BAD NEWS:** there is **no** sparse analog!

[Klep Magron Povh '21]

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**GOOD NEWS:** there is an NC analog of the sparse Putinar's representation!

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[Klep Magron Povh '21]

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# Correlative sparsity

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**Theorem [Klep Magron Povh '21]**

$f \succcurlyeq 0$  on  $\{\underline{x} : s_j(\underline{x}) \succcurlyeq 0\}$

chordal graph  $G$  with cliques  $I_k \Rightarrow$

ball constraints for each  $x(I_k)$

$$f = \sum_{k,i} s_{ki}^* s_{ki} + \sum_{j,i} t_{ji}^* s_j t_{ji}$$

$s_{ki}$  "sees" vars in  $I_k$

$t_{ji}$  "sees" vars from  $s_j$

# Correlative sparsity

---

I<sub>3322</sub> Bell inequality (entanglement in quantum information)

$$f = x_1(y_1 + y_2 + y_3) + x_2(y_1 + y_2 - y_3) + x_3(y_1 - y_2) - x_1 - 2y_1 - y_2$$

Maximal violation levels  $\rightarrow$  **upper bounds** on  $\lambda_{\max}$  of  $f$  on  $\{(x, y) : x_i^2 = x_i, y_j^2 = y_j, x_i y_j = y_j x_i\}$

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6	0.2508753977180	(1 hour)

PERFORMANCE



VS



ACCURACY

# Term sparsity

[Postdoc Wang '19-21] ANR Tremplin-ERC



$$f = x_1^2 - x_1x_2 - x_2x_1 + 3x_2^2 - 2x_1x_2x_1 + 2x_1x_2^2x_1 - x_2x_3 \\ - x_3x_2 + 6x_3^2 + 9x_2^2x_3 + 9x_3^2x_2 - 54x_3x_2x_3 + 142x_3x_2^2x_3$$

[Burgdorf Klep Povh '16] → Newton chip method

$$f = (1 \quad x_1 \quad x_2 \quad x_3 \quad x_2x_1 \quad x_3x_2) \underbrace{Q}_{\neq 0} \begin{pmatrix} 1 \\ x_1 \\ x_2 \\ x_3 \\ x_1x_2 \\ x_2x_3 \end{pmatrix}$$

↪  $\frac{6 \times 7}{2} = 28$  “unknown” entries in  $Q$

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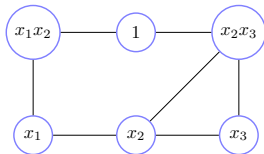
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💡 Term sparsity pattern graph  $G$



# Term sparsity

[Postdoc Wang '19-21] ANR Tremplin-ERC



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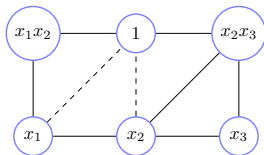
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[Postdoc Wang '19-21] ANR Tremplin-ERC



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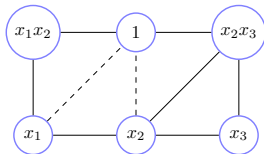
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Replace  $Q$  by  $Q_{G'}$  with nonzero entries at edges of  $G'$

$\rightsquigarrow 6 + 9 = 15$  "unknown" entries in  $Q_{G'}$

# Term sparsity

---

At step  $d$  of NC hierarchy, one considers vector  $\mathbf{W}_d$  of words with length at most  $d$

The **tsp** graph  $G$  has edges  $E$  with

$$\{u, v\} \in E \Leftrightarrow u^*v \in \text{supp}(f) \cup \bigcup_{s \in S} \text{supp}(s) \cup \{w^*w \mid w \in \mathbf{W}_d\}$$

$\rightsquigarrow$  **support extension**

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$\rightsquigarrow$  support extension  $\rightsquigarrow$  chordal extension  $G'$

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$\rightsquigarrow$  **support extension**  $\rightsquigarrow$  **chordal extension**  $G'$

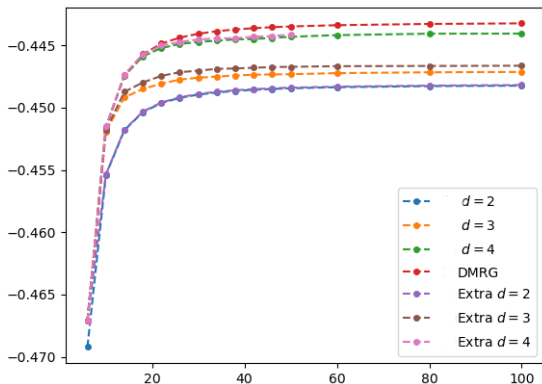
By iteratively performing **support extension** & **chordal extension**

$$G^{(1)} = G' \subseteq \dots \subseteq G^{(\ell)} \subseteq G^{(\ell+1)} \subseteq \dots$$

💡 Two-level hierarchy of lower bounds for  $\lambda_{\min}(f)$  on  $\mathcal{D}_S^\infty$ , indexed by sparse order  $\ell$  and relaxation order  $d$

# An example from condensed matter

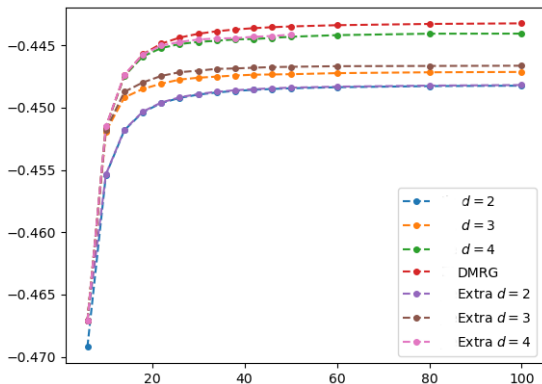
Hamiltonian ground-state energy    1D lattice



**Dense**  $d = 4, n = 10^2 \Rightarrow 10^{19}$  variables (solvers handle  $\simeq 10^4$ )

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**Sparse** solved in  $\simeq 1$  hour on the lab cluster

# Conclusion and perspectives

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SPARSITY EXPLOITING CONVERGING HIERARCHIES to minimize eigenvalue/trace

FAST IMPLEMENTATION IN JULIA: NCTSSOS

💡 Combine correlative & term sparsity  $\rightsquigarrow$  solves problems with 100-1000 variables

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💡 **SYMMETRIC** noncommutative problems?  
Ground state energy of hamiltonians    💡 **symmetric & sparse**

# Thank you for your attention!

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<https://homepages.laas.fr/vmagron>



Wang & Magron. Exploiting term sparsity in Noncommutative Polynomial Optimization. *Computational Optimization & Applications*, arxiv:2010.06956

[NCTSSOS](#)



Klep, Magron & Povh. Sparse Noncommutative Polynomial Optimization. *Math Prog. A*, arxiv:1909.00569

[NCSOSTools](#)