

# Enclosures of Roundoff Errors using SDP

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**Metalibm workshop:**

“Elementary functions, digital filters and beyond”



12-13 March 2018



# Errors and Proofs

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## GUARANTEED OPTIMIZATION



Input : Linear problem  (LP), geometric, semidefinite  (SDP)

Output : solution + **certificate**  numeric-symbolic  $\rightsquigarrow$   formal

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## VERIFICATION OF CRITICAL SYSTEMS

Reliable software/hardware embedded codes

Aerospace control



molecular biology, robotics, code synthesis, ...



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## Efficient Verification of Nonlinear Systems

- Automated precision tuning of systems/programs  
analysis/synthesis
- Efficiency sparsity correlation patterns
- Certified approximation algorithms

# Roundoff Error Bounds

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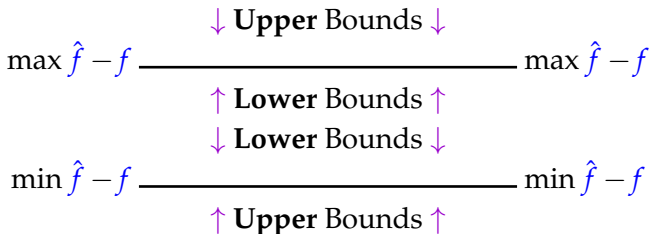
**Real** :  $f(\mathbf{x}) := x_1 \times x_2 + x_3$

**Floating-point** :  $\hat{f}(\mathbf{x}, \mathbf{e}) := [x_1 x_2 (1 + e_1) + x_3](1 + e_2)$

Input variable constraints  $\mathbf{x} \in \mathbf{X}$

Finite precision  $\rightsquigarrow$  bounds over  $\mathbf{e} \in \mathbf{E}$ :  $|e_i| \leq 2^{-53}$  (double)

**Guarantees** on absolute round-off error  $|\hat{f} - f|$  ?



# Nonlinear Programs

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- **Polynomials** programs : +, -, ×

$$x_2x_5 + x_3x_6 + x_1(-x_1 + x_2 + x_3 - x_4 + x_5 + x_6)$$

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- **Semialgebraic** programs: |·|, √, /, sup, inf

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- **Transcendental** programs: arctan, exp, log, ...

$$\log(1 + \exp(x))$$



# Existing Frameworks

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## Classical methods:

- Abstract domains [Goubault-Putot 11]

FLUCTUAT: intervals, octagons, zonotopes

- Interval arithmetic [Daumas-Melquiond 10]

GAPPA: interface with COQ proof assistant

# Existing Frameworks

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## Recent progress:

- Affine arithmetic + SMT

[Darulova 14], [Darulova 17], [Izycheva 17]

rosa: sound compiler for reals (SCALA)

Daisy: roundoff error + rewriting + multi-precision

- Symbolic Taylor expansions [Solovyev 15], [Chiang 17]

FPTaylor: certified optimization (OCAML/HOL-LIGHT)

- Improve numerical accuracy of programs [Damouche 17]

Salsa: program transformation

- Guided random testing s3fp [Chiang 14]

# Contributions

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Maximal Roundoff error of the program implementation of  $f$ :

$$r^* := \max |\hat{f}(\mathbf{x}, \mathbf{e}) - f(\mathbf{x})|$$

**Decomposition:** **linear** term  $l$  w.r.t.  $\mathbf{e}$  + nonlinear term  $h$

$$\max |l| + \max |h| \geq r^* \geq \max |l| - \max |h|$$

- Coarse bound of  $h$  with interval arithmetic
- **Semidefinite programming** (SDP) bounds for  $l$ :

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↓ Upper Bounds ↓

↑ Lower Bounds ↑

↓ Lower Bounds ↓

↑ Upper Bounds ↑

**Sparse SDP relaxations**

**Robust SDP relaxations**

# Contributions

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- 1 General **SDP** framework for **upper** and **lower** bounds
- 2 **Comparison** with SMT & affine/Taylor arithmetic:  
     $\rightsquigarrow$  **Efficient** optimization  $\oplus$  **Tight** upper bounds
- 3 Extensions to **transcendental**/conditional programs
- 4 Formal verification of SDP bounds 🐔
- 5 Open source tools:

Real2Float (in OCAML and COQ)

↓ Upper Bounds ↓

↑ Upper Bounds ↑

FPSPD (in MATLAB)

↑ Lower Bounds ↑

↓ Lower Bounds ↓

Introduction

**Semidefinite Programming for Polynomial Optimization**

Upper Bounds with Sparse SDP

Lower Bounds with Robust SDP

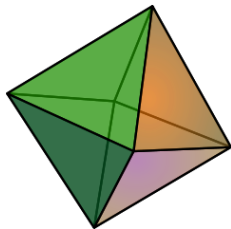
Conclusion

# What is Semidefinite Programming?

---

- Linear Programming (LP):

$$\begin{aligned} \min_{\mathbf{z}} \quad & \mathbf{c}^\top \mathbf{z} \\ \text{s.t.} \quad & \mathbf{A} \mathbf{z} \geq \mathbf{d} . \end{aligned}$$



- Linear cost  $\mathbf{c}$
- Linear inequalities “ $\sum_i A_{ij} z_j \geq d_i$ ”

**Polyhedron**

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- Linear cost  $\mathbf{c}$
- Symmetric matrices  $\mathbf{F}_0, \mathbf{F}_i$
- Linear matrix inequalities “ $\mathbf{F} \succcurlyeq 0$ ”  
( $\mathbf{F}$  has nonnegative eigenvalues)



**Spectrahedron**



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Spectrahedron

# SDP for Polynomial Optimization

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- Prove **polynomial inequalities** with SDP:

$$f(a, b) := a^2 - 2ab + b^2 \geq 0 .$$

- Find  $\mathbf{z}$  s.t.  $f(a, b) = \underbrace{\begin{pmatrix} a & b \\ z_1 & z_2 \\ z_2 & z_3 \end{pmatrix}}_{\succcurlyeq 0} \begin{pmatrix} a \\ b \end{pmatrix}$ .

- Find  $\mathbf{z}$  s.t.  $a^2 - 2ab + b^2 = z_1 a^2 + 2z_2 ab + z_3 b^2 \quad (\mathbf{A} \mathbf{z} = \mathbf{d})$

$$\begin{pmatrix} z_1 & z_2 \\ z_2 & z_3 \end{pmatrix} = \underbrace{\begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix}}_{\mathbf{F}_1} z_1 + \underbrace{\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}}_{\mathbf{F}_2} z_2 + \underbrace{\begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}}_{\mathbf{F}_3} z_3 \succcurlyeq \underbrace{\begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix}}_{\mathbf{F}_0}$$

# SDP for Polynomial Optimization

---

- Choose a cost  $\mathbf{c}$  e.g.  $(1, 0, 1)$  and solve:

$$\begin{aligned} \min_{\mathbf{z}} \quad & \mathbf{c}^\top \mathbf{z} \\ \text{s.t.} \quad & \sum_i \mathbf{F}_i z_i \succcurlyeq \mathbf{F}_0, \quad \mathbf{A} \mathbf{z} = \mathbf{d}. \end{aligned}$$

- Solution  $\begin{pmatrix} z_1 & z_2 \\ z_2 & z_3 \end{pmatrix} = \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix} \succcurlyeq 0$  (eigenvalues 0 and 2)

- $a^2 - 2ab + b^2 = (a \ b) \underbrace{\begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix}}_{\succcurlyeq 0} \begin{pmatrix} a \\ b \end{pmatrix} = (a - b)^2.$

- Solving **SDP**  $\implies$  Finding **SUMS OF SQUARES** certificates

# SDP for Polynomial Optimization

---

NP hard General Problem:  $f^* := \min_{\mathbf{x} \in \mathbf{X}} f(\mathbf{x})$

- Semialgebraic set  $\mathbf{X} := \{\mathbf{x} \in \mathbb{R}^n : g_1(\mathbf{x}) \geq 0, \dots, g_m(\mathbf{x}) \geq 0\}$

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$$\underbrace{x_1 x_2}_f + \frac{1}{8} = \frac{1}{2} \overbrace{\left(x_1 + x_2 - \frac{1}{2}\right)^2}^{\sigma_0} + \frac{1}{2} \overbrace{x_1(1 - x_1)}^{\sigma_1} + \frac{1}{2} \overbrace{x_2(1 - x_2)}^{\sigma_2}$$

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- Sums of squares (SOS)  $\sigma_i$

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- Sums of squares (SOS)  $\sigma_i$
- Bounded degree:

$$\mathcal{Q}_k(\mathbf{X}) := \left\{ \sigma_0 + \sum_{j=1}^m \sigma_j g_j, \text{ with } \deg \sigma_j g_j \leq 2k \right\}$$



# SDP for Polynomial Optimization

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- **Hierarchy of SDP relaxations:**

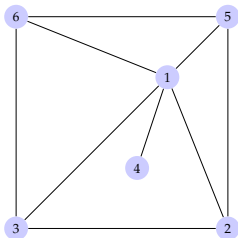
$$\lambda_k := \sup_{\lambda} \left\{ \lambda : f - \lambda \in \mathcal{Q}_k(\mathbf{X}) \right\}$$

- Convergence guarantees  $\lambda_k \uparrow f^*$  [Lasserre 01]
- Can be computed with SDP solvers (CSDP, SDPA)
- **“No Free Lunch” Rule:**  $\binom{n+2k}{n}$  SDP variables
- Extension to **semialgebraic functions**  
 $\rightsquigarrow r(\mathbf{x}) = p(\mathbf{x}) / \sqrt{q(\mathbf{x})}$  [Lasserre-Putinar 10]

# Sparse SDP Optimization [Waki, Lasserre 06]

- Correlative sparsity pattern (csp) of variables

$$x_2x_5 + x_3x_6 - x_2x_3 - x_5x_6 + x_1(-x_1 + x_2 + x_3 - x_4 + x_5 + x_6)$$



1 Maximal cliques  $C_1, \dots, C_l$

2 Average size  $\kappa \rightsquigarrow \binom{\kappa+2k}{\kappa}$   
variables

$$C_1 := \{1, 4\}$$

$$C_2 := \{1, 2, 3, 5\}$$

$$C_3 := \{1, 3, 5, 6\}$$

Dense SDP: 210 variables

Sparse SDP: 115 variables

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# Polynomial Programs

**Input:** exact  $f(\mathbf{x})$ , approx  $\hat{f}(\mathbf{x}, \mathbf{e})$ ,  $\mathbf{x} \in \mathbf{X}$ ,  $|e_i| \leq 2^{-53}$

**Output:** Bound for  $f - \hat{f}$

- 1: Error  $r(\mathbf{x}, \mathbf{e}) := f(\mathbf{x}) - \hat{f}(\mathbf{x}, \mathbf{e}) = \sum_{\alpha} r_{\alpha}(\mathbf{e}) \mathbf{x}^{\alpha}$

- 2: Decompose  $r(\mathbf{x}, \mathbf{e}) = l(\mathbf{x}, \mathbf{e}) + h(\mathbf{x}, \mathbf{e})$

- 3:  $l(\mathbf{x}, \mathbf{e}) = \sum_{i=1}^m s_i(\mathbf{x}) e_i$

- 4: Maximal cliques correspond to  $\{\mathbf{x}, e_1\}, \dots, \{\mathbf{x}, e_m\}$

**Dense** relaxation:  $\binom{n+m+2k}{n+m}$  SDP variables

**Sparse** relaxation:  $m \binom{n+1+2k}{n+1}$  SDP variables

# Preliminary Comparisons

$$f(\mathbf{x}) := x_2x_5 + x_3x_6 - x_2x_3 - x_5x_6 + x_1(-x_1 + x_2 + x_3 - x_4 + x_5 + x_6)$$

$$\mathbf{x} \in [4.00, 6.36]^6, \quad \mathbf{e} \in [-\epsilon, \epsilon]^{15}, \quad \epsilon = 2^{-53}$$

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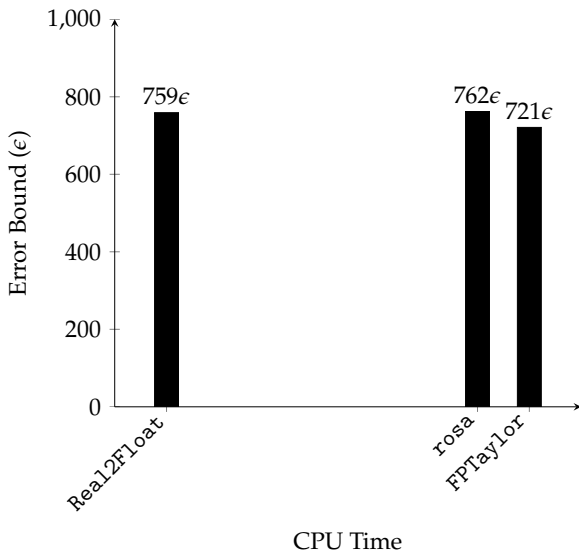
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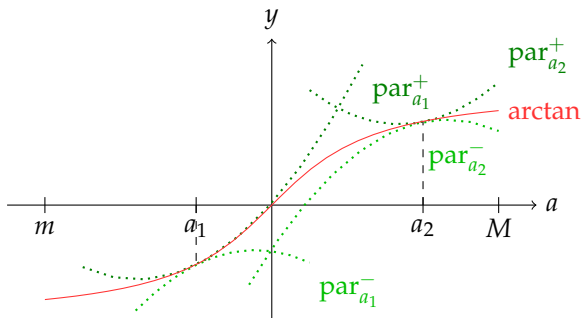
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- **Symbolic Taylor** FPTaylor tool:  $721\epsilon$  ( $21 \times$  more CPU)
- **SMT-based** rosa tool:  $762\epsilon$  ( $19 \times$  more CPU)

# Preliminary Comparisons



# Extensions: Transcendental Programs

Reduce  $f^* := \inf_{\mathbf{x} \in \mathbf{K}} f(\mathbf{x})$  to semialgebraic optimization



# Extensions: Conditionals

if ( $p(\mathbf{x}) \leq 0$ )  $f(\mathbf{x})$ ; else  $g(\mathbf{x})$ ;

DIVERGENCE PATH ERROR:

$$r^* := \max\left\{ \begin{array}{l} \max_{p(\mathbf{x}) \leq 0, p(\mathbf{x}, \mathbf{e}) \geq 0} |\hat{f}(\mathbf{x}, \mathbf{e}) - g(\mathbf{x})| \\ \max_{p(\mathbf{x}) \geq 0, p(\mathbf{x}, \mathbf{e}) \leq 0} |\hat{g}(\mathbf{x}, \mathbf{e}) - f(\mathbf{x})| \\ \max_{p(\mathbf{x}) \geq 0, p(\mathbf{x}, \mathbf{e}) \geq 0} |\hat{f}(\mathbf{x}, \mathbf{e}) - f(\mathbf{x})| \\ \max_{p(\mathbf{x}) \leq 0, p(\mathbf{x}, \mathbf{e}) \leq 0} |\hat{g}(\mathbf{x}, \mathbf{e}) - g(\mathbf{x})| \end{array} \right\}$$

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**Lower Bounds with Robust SDP**

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# Method 1: geneig [Lasserre 11]

Generalized eigenvalue problem:

$$f^* := \min_{\mathbf{x} \in \mathbf{X}} f(\mathbf{x}) \leq \lambda_k := \sup_{\lambda} \lambda$$

s.t.  $\mathbf{M}_k(f \mathbf{y}) \succeq \lambda \mathbf{M}_k(\mathbf{y})$ .

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Uniform distribution moments:  $\mathbf{y}_\alpha := \int_{\mathbf{X}} \mathbf{x}^\alpha d\mathbf{x}$

Localizing matrices  $\mathbf{M}_k(f \mathbf{y})$ :

$$\mathbf{M}_1(f \mathbf{y}) = \begin{matrix} & \begin{matrix} 1 & x_1 & x_2 \end{matrix} \\ \begin{matrix} 1 \\ x_1 \\ x_2 \end{matrix} & \begin{pmatrix} \int_{\mathbf{X}} f(\mathbf{x}) d\mathbf{x} & \int_{\mathbf{X}} f(\mathbf{x}) x_1 d\mathbf{x} & \int_{\mathbf{X}} f(\mathbf{x}) x_2 d\mathbf{x} \\ \int_{\mathbf{X}} f(\mathbf{x}) x_1 d\mathbf{x} & \int_{\mathbf{X}} f(\mathbf{x}) x_1^2 d\mathbf{x} & \int_{\mathbf{X}} f(\mathbf{x}) x_1 x_2 d\mathbf{x} \\ \int_{\mathbf{X}} f(\mathbf{x}) x_2 d\mathbf{x} & \int_{\mathbf{X}} f(\mathbf{x}) x_2 x_1 d\mathbf{x} & \int_{\mathbf{X}} f(\mathbf{x}) x_2^2 d\mathbf{x} \end{pmatrix} \end{matrix}$$

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**Theorem [Lasserre 11, de Klerk et al. 17]**

$$\lambda_k \downarrow f^* \quad \text{and} \quad \lambda_k - f^* = \mathcal{O}(1/\sqrt{k})$$



Elementary calculation with  $f(\mathbf{x}) = \sum_{\alpha} f_{\alpha} \mathbf{x}^{\alpha}$ :

$$f^* := \min_{\mathbf{x} \in \mathbf{X}} f(\mathbf{x}) \leq f_k^H := \min_{|\eta + \beta| \leq 2k} \sum_{\alpha} f_{\alpha} \frac{\gamma_{\eta + \alpha, \beta}}{\gamma_{\eta, \beta}}$$

## Method 2: mvbeta [DeKlerk et al. 17]

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Multivariate beta distribution moments:

$$\gamma_{\eta, \beta} := \int_{\mathbf{X}} \mathbf{x}^{\eta} (1 - \mathbf{x})^{\beta} d\mathbf{x}.$$

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$$\gamma_{\eta, \beta} := \int_{\mathbf{X}} \mathbf{x}^{\eta} (1 - \mathbf{x})^{\beta} d\mathbf{x}.$$

Theorem [DeKlerk et al. 17]

$$f_k^H \downarrow f^* \quad \text{and} \quad \lambda_k - f^* = \mathcal{O}(1/\sqrt{k})$$

## Method 3: robustsdp

Robust SDP with  $l(\mathbf{x}, \mathbf{e}) = \sum_{i=1}^m s_i(\mathbf{x})e_i$ :

$$l^* := \min_{(\mathbf{x}, \mathbf{e}) \in \mathbf{X} \times \mathbf{E}} l(\mathbf{x}, \mathbf{e}) \leq \lambda'_k := \sup_{\lambda} \lambda$$

$$\text{s.t. } \forall \mathbf{e} \in \mathbf{E}, \mathbf{M}_k(l \mathbf{y}) \succeq \lambda \mathbf{M}_k(\mathbf{y}).$$

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$$\text{Linearity} \rightsquigarrow \mathbf{M}_k(l \mathbf{y}) = \sum_{i=1}^m e_i \mathbf{M}_k(s_i \mathbf{y})$$

$$\text{Factorization} \rightsquigarrow \mathbf{M}_k(s_i \mathbf{y}) = \mathbf{L}_k^i \mathbf{R}_k^i$$

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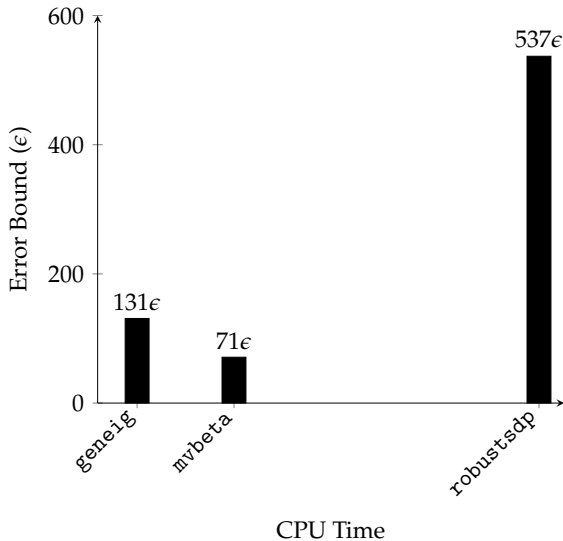
$$\text{Factorization } \rightsquigarrow \mathbf{M}_k(s_i \mathbf{y}) = \mathbf{L}_k^i \mathbf{R}_k^i$$

Theorem following from [El Ghaoui et al. 98]

$$\lambda'_k \downarrow l^* \text{ and } \lambda'_k = \sup_{\lambda, \mathbf{S}, \mathbf{G}} \lambda$$

$$\text{s.t. } \begin{pmatrix} -\lambda \mathbf{M}_k(\mathbf{y}) - \mathbf{L}_k \mathbf{S} \mathbf{L}_k^T & \mathbf{R}_k^T + \mathbf{L}_k \mathbf{G} \\ \mathbf{R}_k - \mathbf{G} \mathbf{L}_k^T & \mathbf{S} \end{pmatrix} \succeq 0, \\ \mathbf{S}^T = \mathbf{S}, \mathbf{G}^T = -\mathbf{G}.$$

# Benchmark kepler0 with $k = 2$



Introduction

Semidefinite Programming for Polynomial Optimization

Upper Bounds with Sparse SDP

Lower Bounds with Robust SDP

**Conclusion**





# Conclusion

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## Sparse/Robust SDP relaxations for NONLINEAR PROGRAMS:

- Polynomial and transcendental programs

- Certified   $\rightsquigarrow$  Formal  roundoff error bounds  
(Joint work with T. Weisser and B. Werner)

- Real2Float and FPSDP open source tools:

<http://nl-certify.forge.ocamlcore.org/real2float.html>

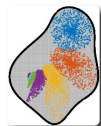
<https://github.com/magronv/FPSDP>

# Conclusion

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## Further research:

- Automatic **FPGA** code generation
- Handling `while/for` loops



**Master / PhD Positions Available !**



# End

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Thank you for your attention!

<http://www-verimag.imag.fr/~magron>

- V. Magron, G. Constantinides, A. Donaldson. Certified Roundoff Error Bounds Using Semidefinite Programming, *ACM Trans. Math. Softw.*, 2017.

↓ Upper Bounds ↓

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↑ Upper Bounds ↑

- V. Magron. Interval Enclosures of Upper Bounds of Roundoff Errors using Semidefinite Programming, [arxiv.org/abs/1611.01318](http://arxiv.org/abs/1611.01318).

↑ Lower Bounds ↑

↓ Lower Bounds ↓

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